

ESc101 : Fundamental of Computing

I Semester 2008-09

Lecture 30

- String (Comparisons)
- Details of method invocation
- Recursion

Comparing two String objects

`str.compareTo(anotherString)`

returns **int** which is less than or equal to or greater than 0 if the String `str` is less than or equal to or greater than string `anotherString`.

Note that the strings are compared lexicographically (like in dictionary). So

- “Z” is *greater* “ABC”
- “Y” is *less* than “YA”;
- “abc” is *greater* than “aabcdef”

What if there are special characters ?

Comparing two string objects

General algorithm :

Let **s** and **t** are two strings to be compared.

- start scanning **s** and **t** from left to right character by character until you find a mismatch. The order (less than or greater than) between the two strings is determined by the comparison between the **unicode value** of the two characters.

Consider the following program

```
public static int g(int j)
{   int i; i= j*2;
    return i;
}
public static int f(int i)
{   int j = g(i)*i;
    return j;
}
public static void main(String args[])
{   double d; int k=2;
    int i = f(k);
    System.out.println(i);
    ...
}
```

what is the output ?

Consider the following program

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public static int g(int j)
{   int i; i= j*2;
    return i;
}
public static int f(int i)
{   int j = g(i)*i;
    return j;
}
public static void main(String args[])
{   double d; int k=2;
    int i = f(k);
    System.out.println(i);
    ...
}
```

output = **8**

What about the following program ?

```
public static int f(int i)
{ if(i<=1)  return 1;
  else return f(i-1)*i;
}
public static void main(String args[])
{  int i = 3;
   int j = f(i);
   System.out.println(j);
}
```

What about the following program ?

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public static int f(int i)
{ if(i<=1)  return 1;
  else return f(i-1)*i;
}
public static void main(String args[])
{  int i = 3;
   int j = f(i);
   System.out.println(j);
}
```

Not clear ??

We need to know the invocation of methods in detail

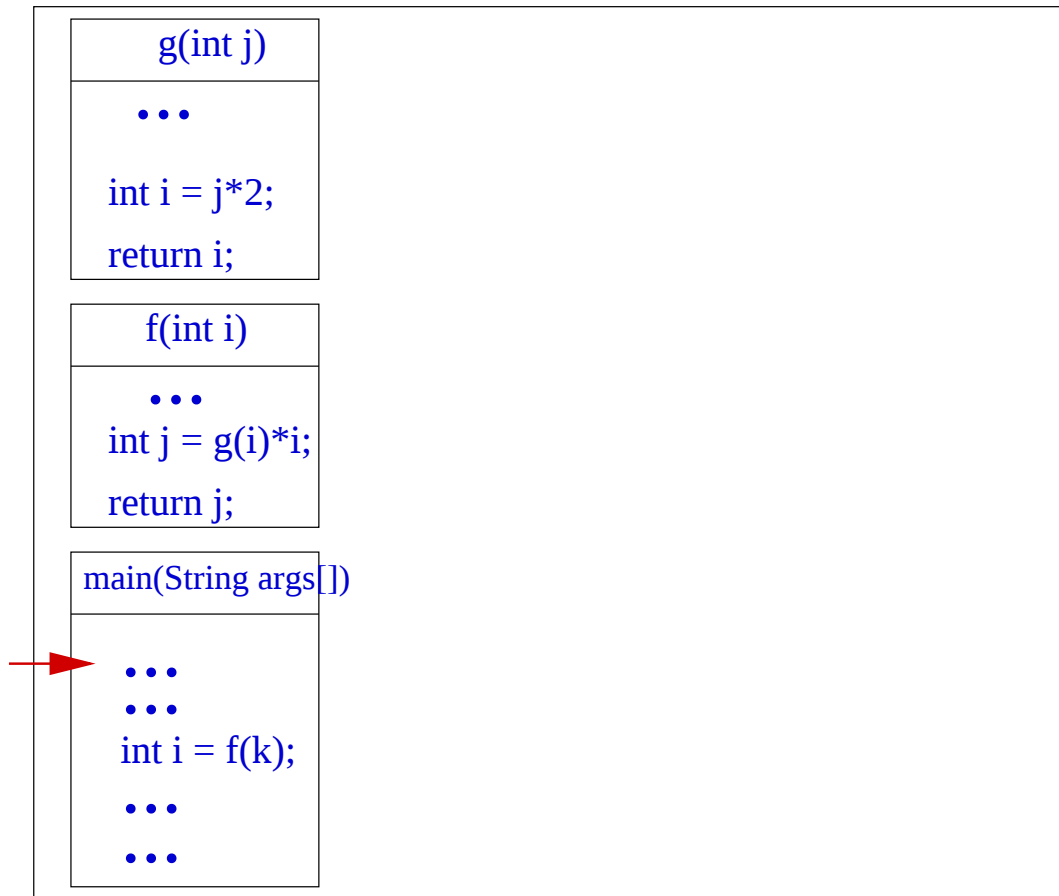
Suppose **M1** invokes **M2**

- How do we keep track of the scope of variables of M1 and M2 at any stage?
- How is it determined where to go at the end of execution of a M2 ?
- How is it ensured that on return from a method **M2** we are in the same state of **M1** in which we left **M1** ?

Let us consider the execution of the first program

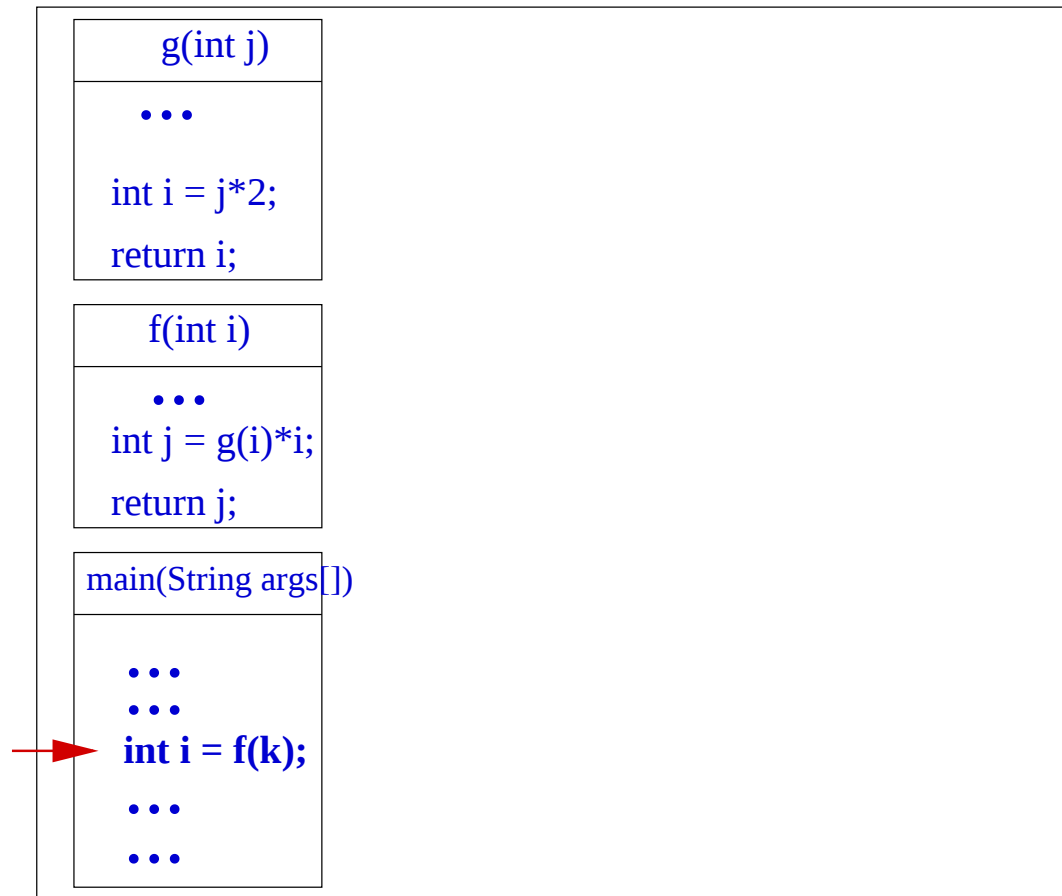
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    return j;
}
public static void main(String args[])
{   double d; int k=2;
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    ...
}
```

Execution starts from **main**



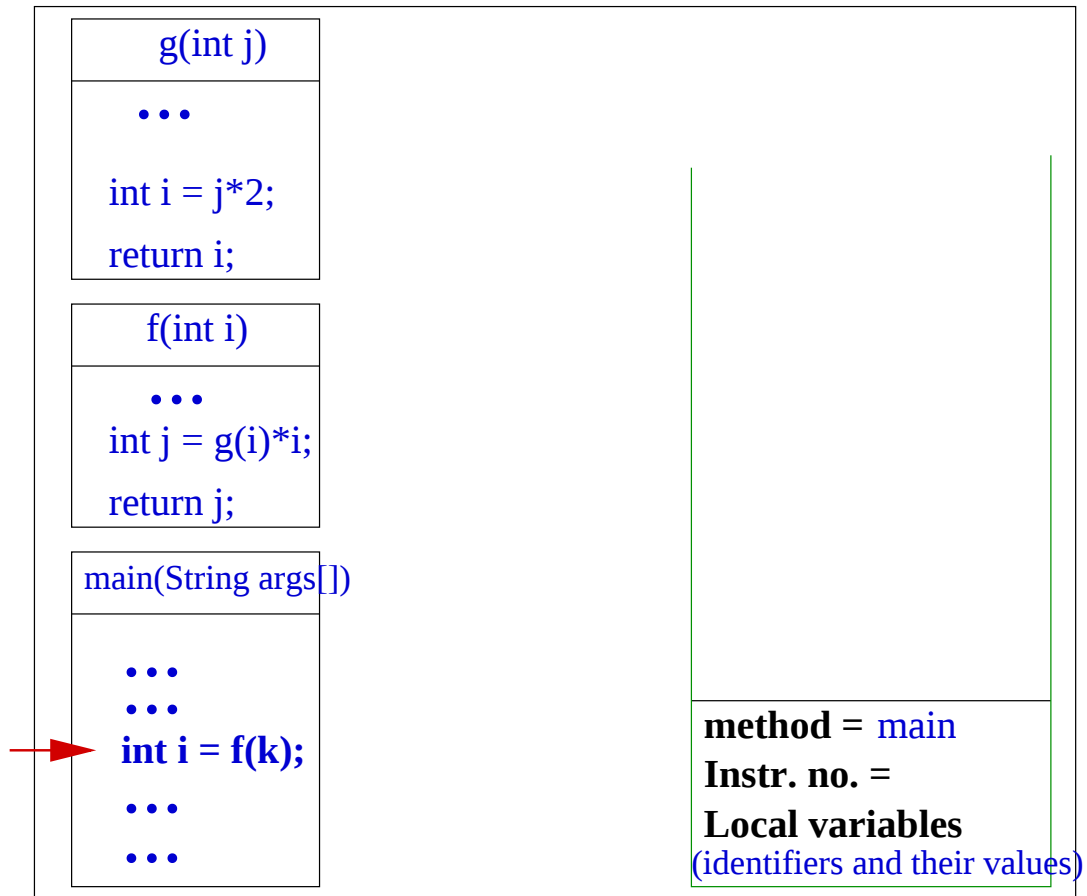
Memory

What actually happens on reaching `int i = f(k)`



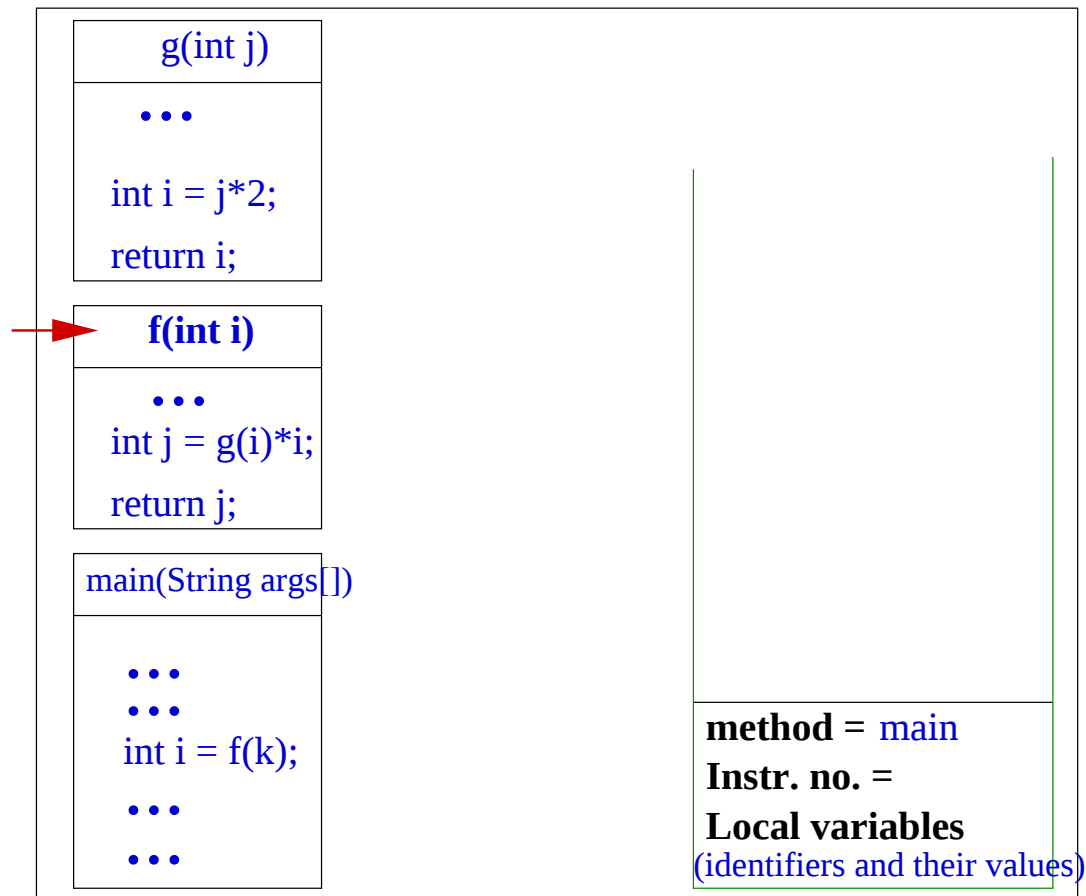
Memory

We keep some records



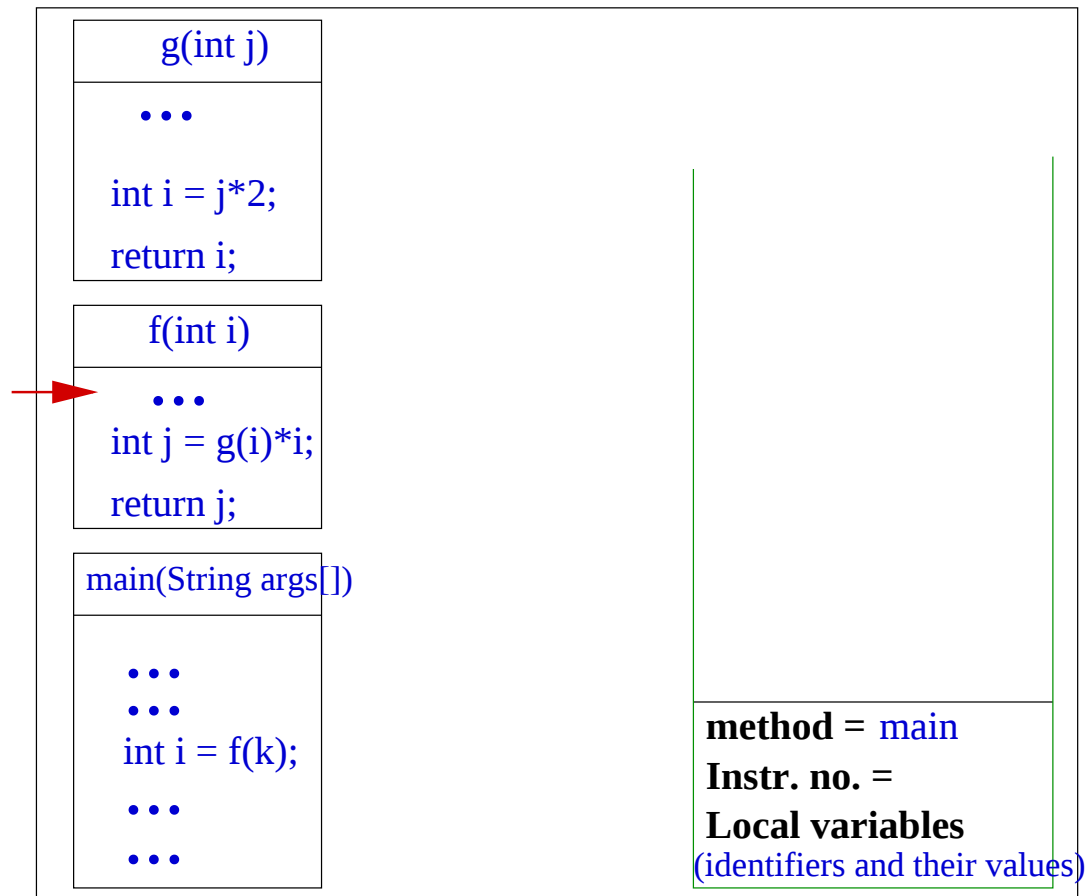
Memory

Then we start excuting **f()**



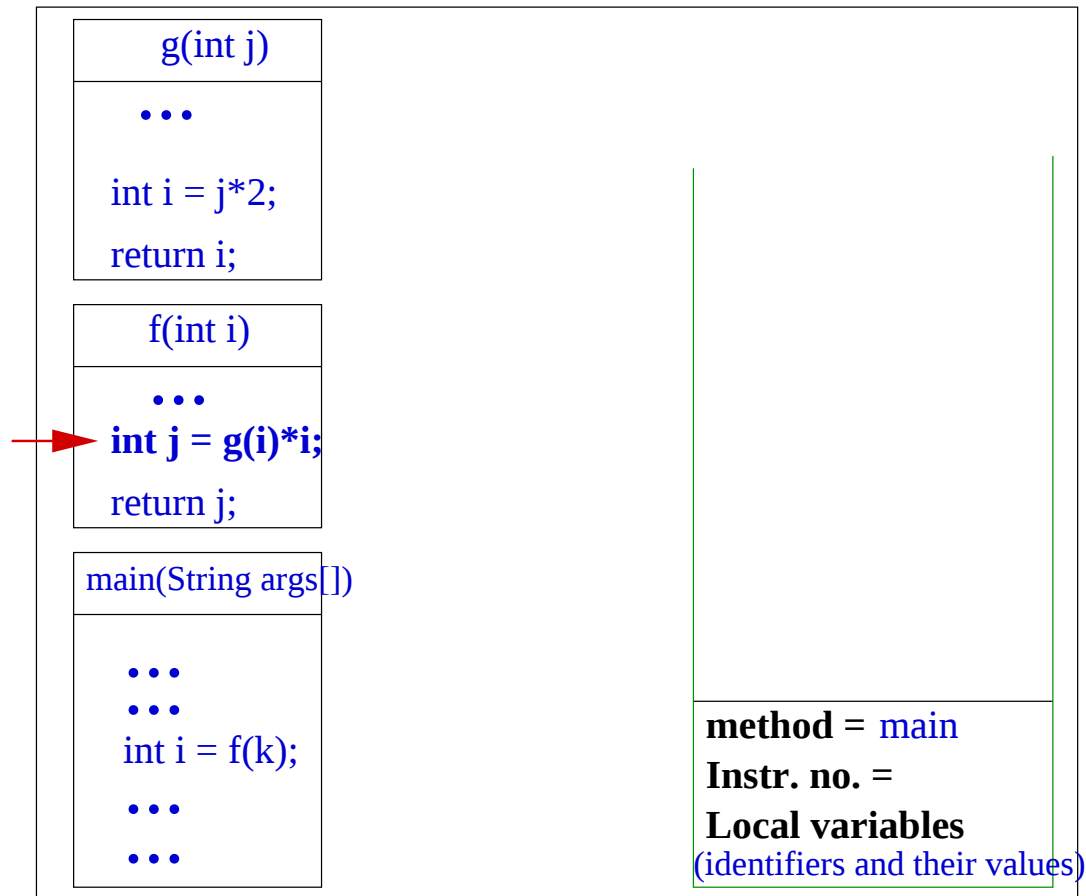
Memory

Execution of **f ()** proceeds from instr. #1



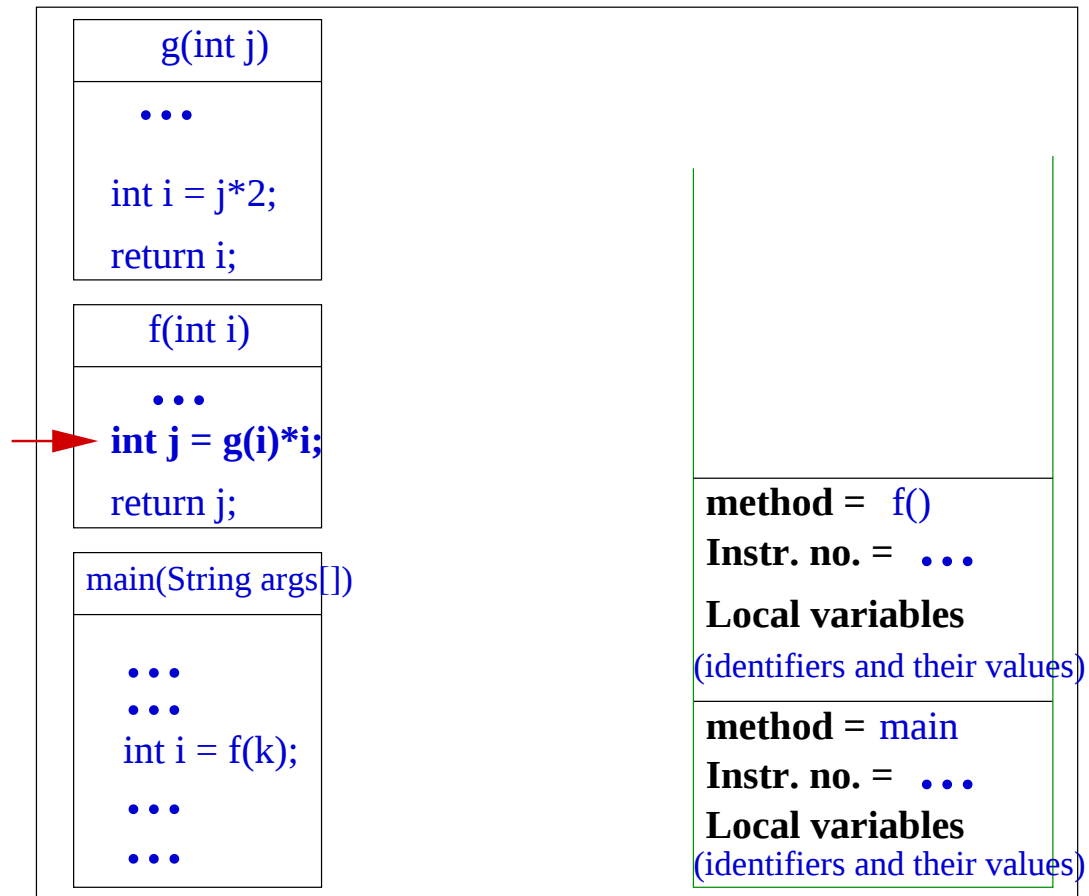
Memory

What happens on reaching `int j = g(i)*i` ?



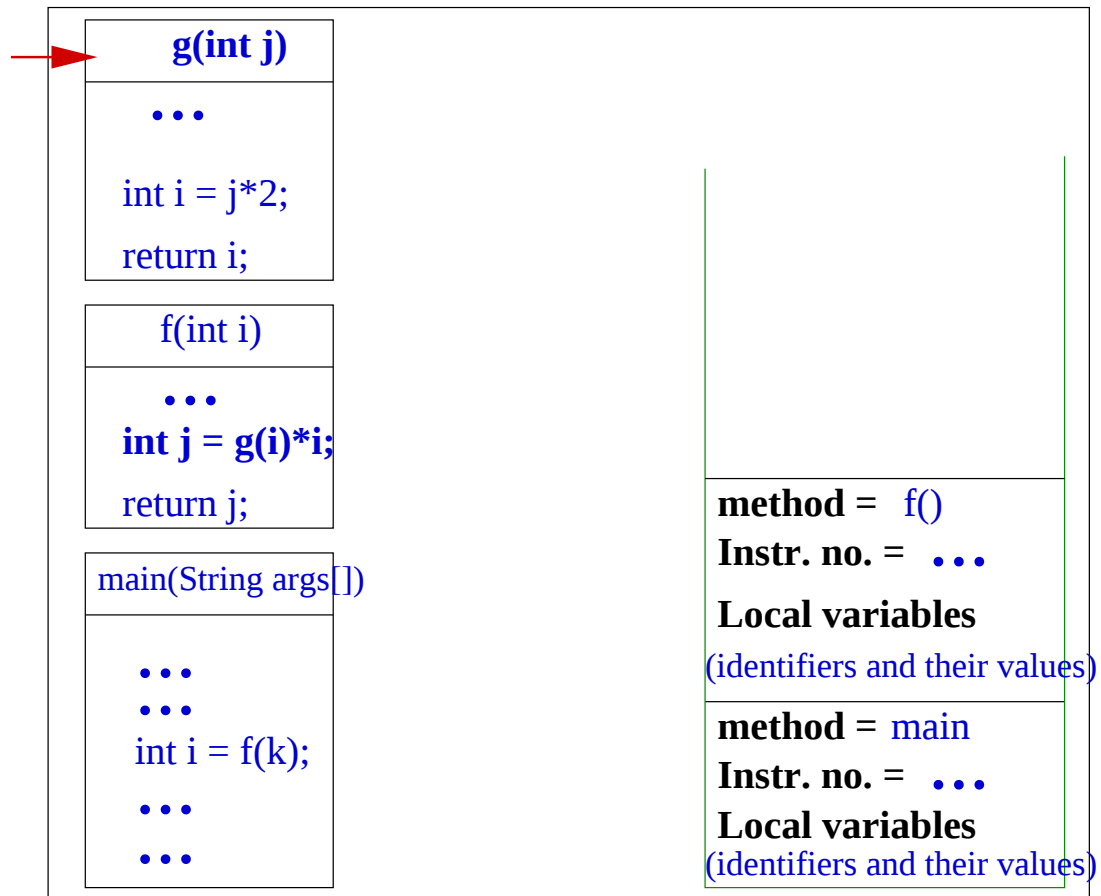
Memory

We keep few more records on the stack

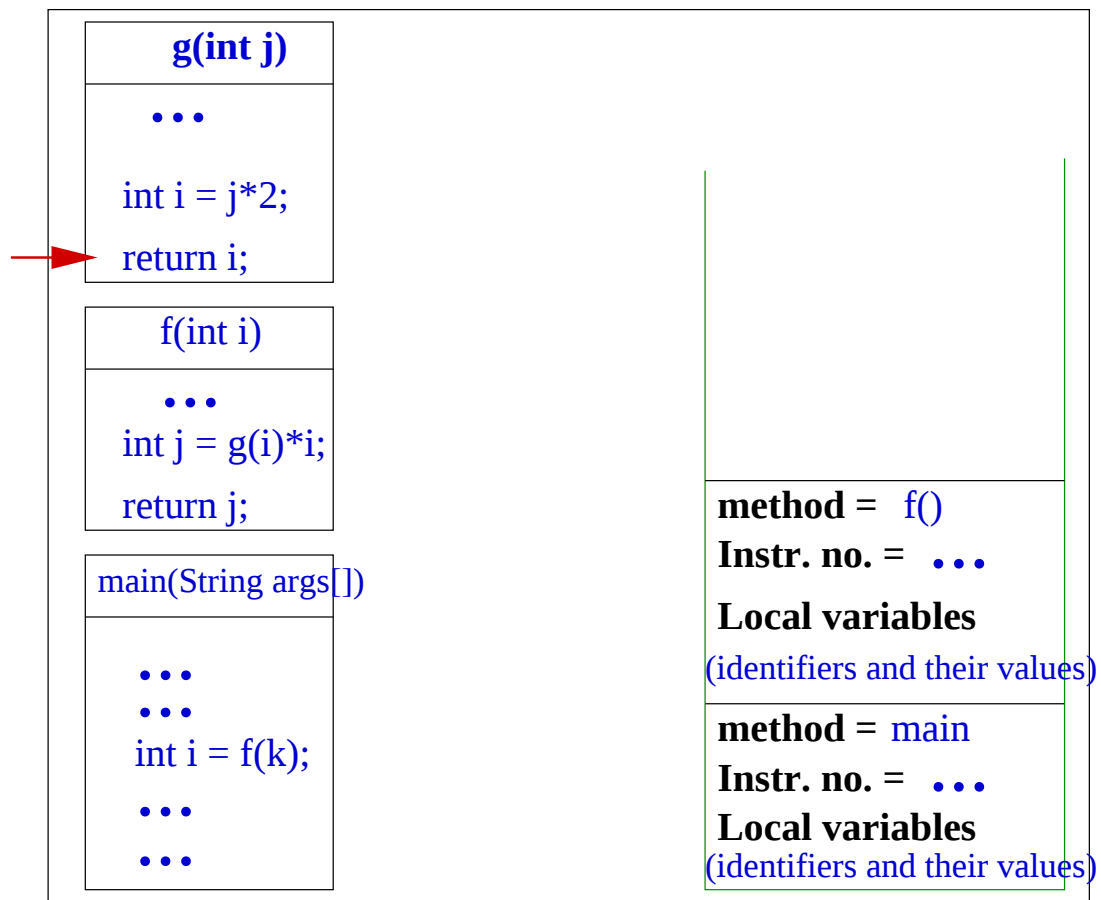


Memory

We start executing `g()`

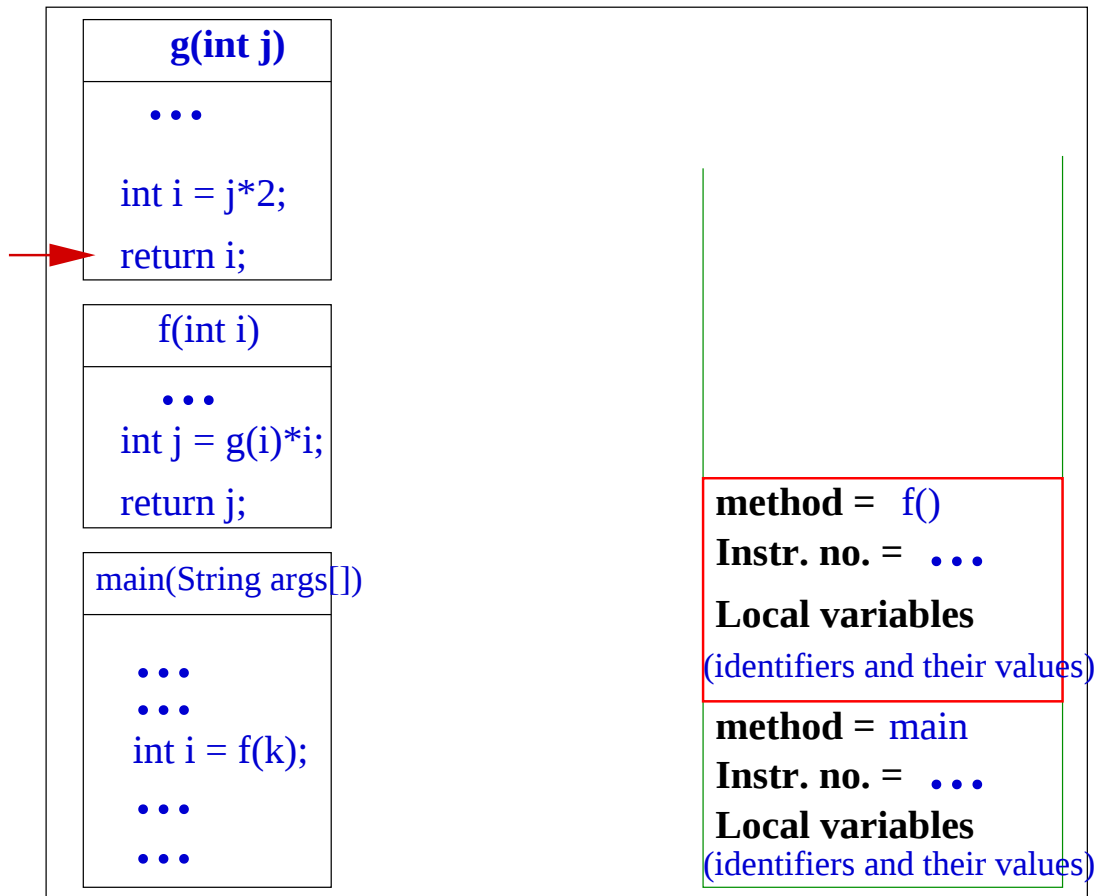


Memory



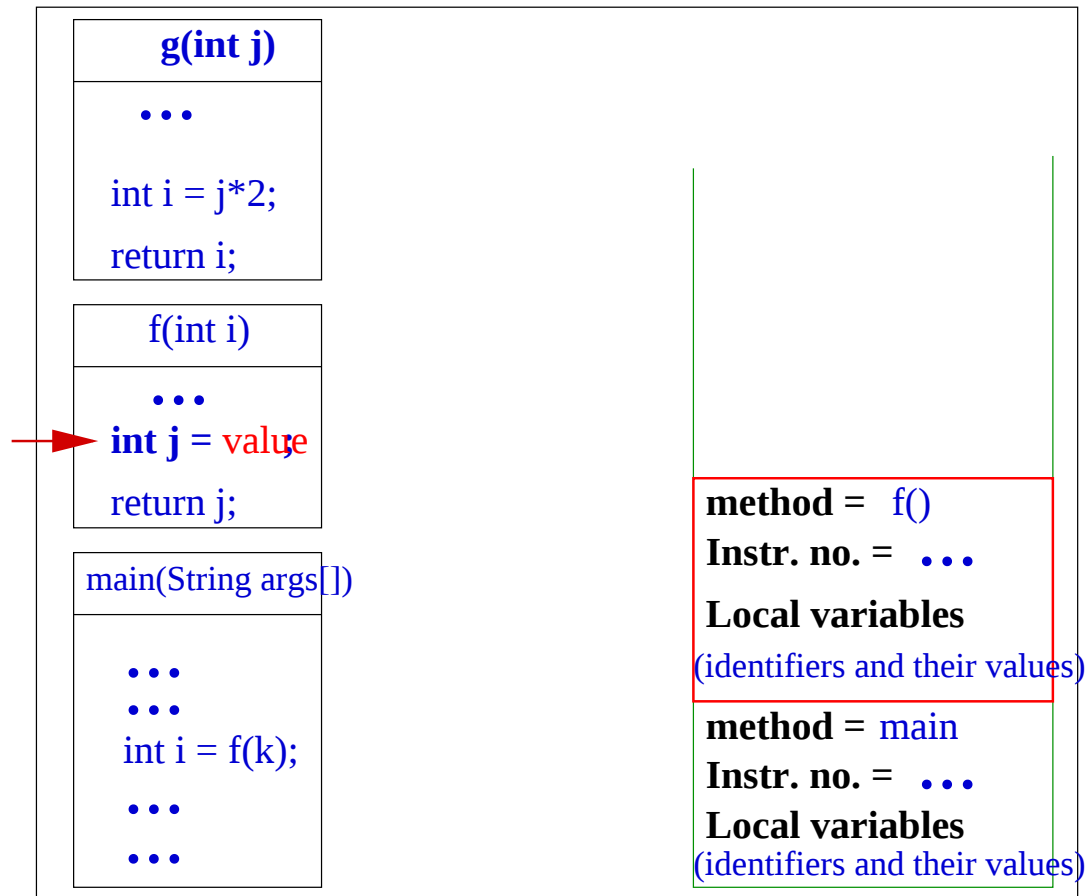
Memory

What happens when we reach end of **g()** ?



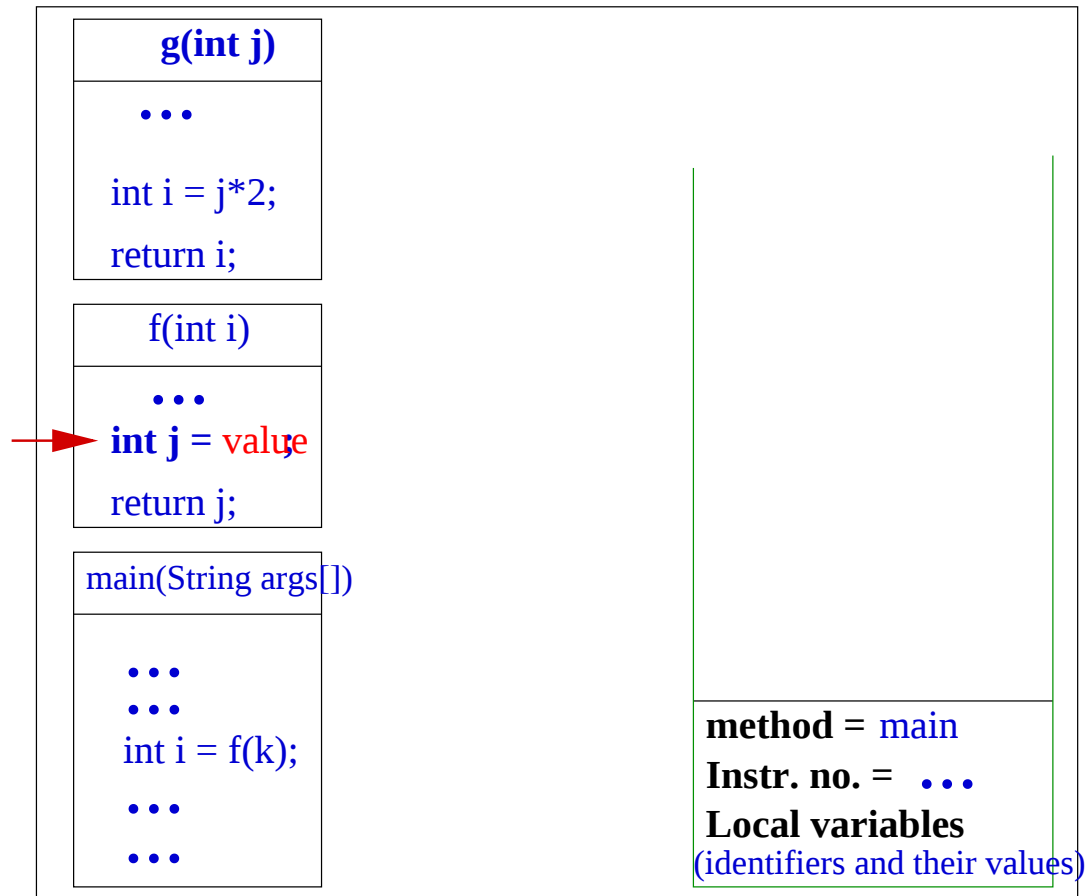
Memory

We explore the top entry of stack



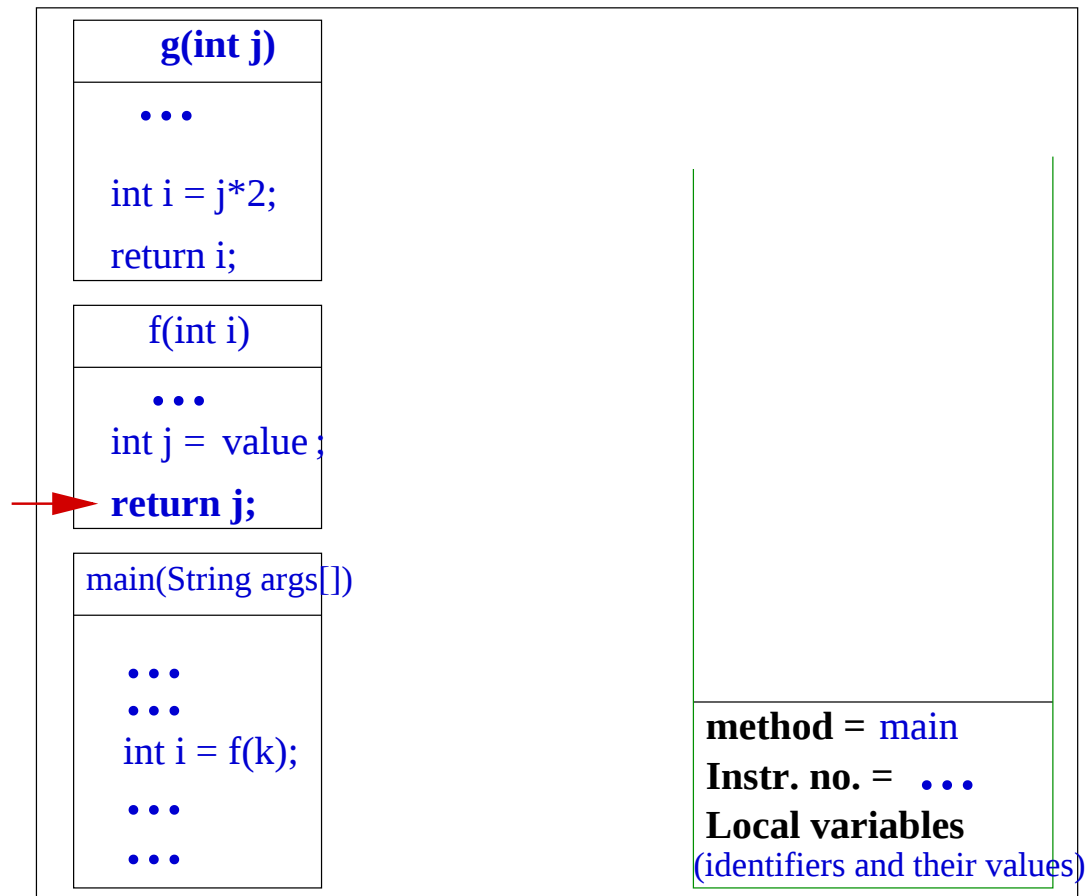
Memory

We copy the value returned by g to j



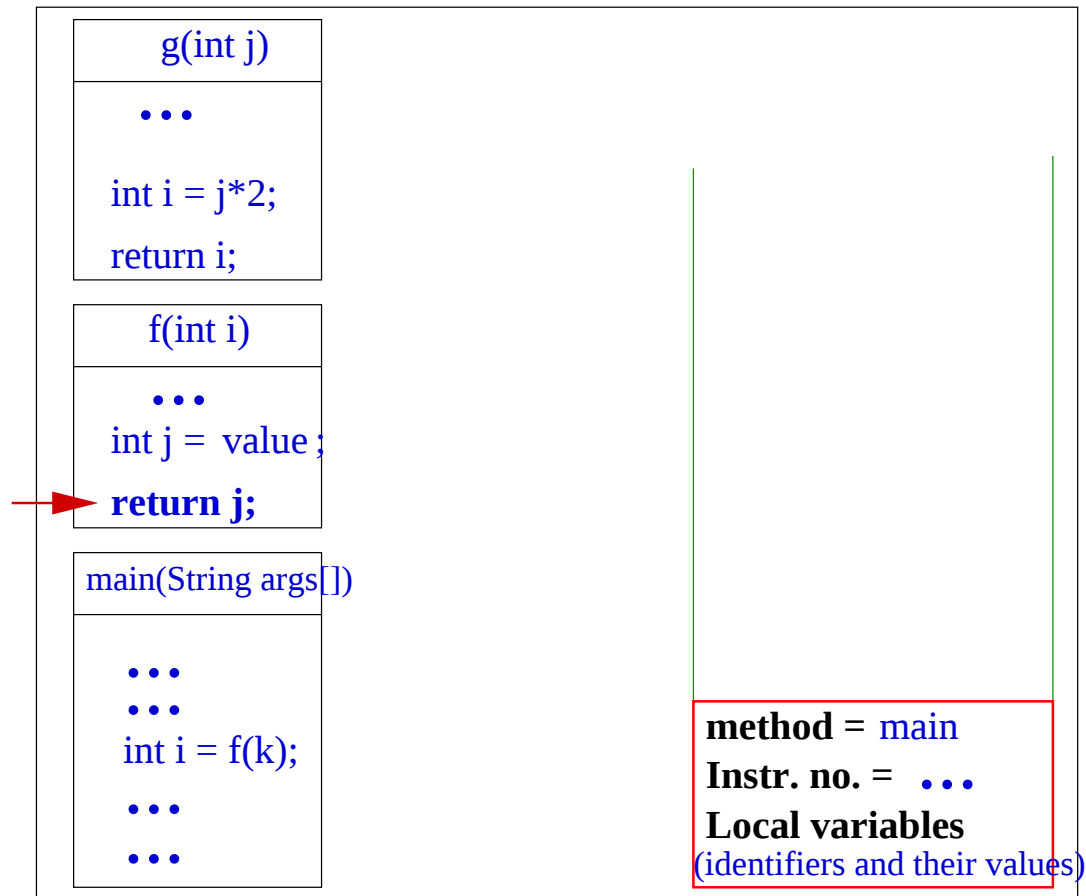
Memory

What happens when we reach the end of **f** ?



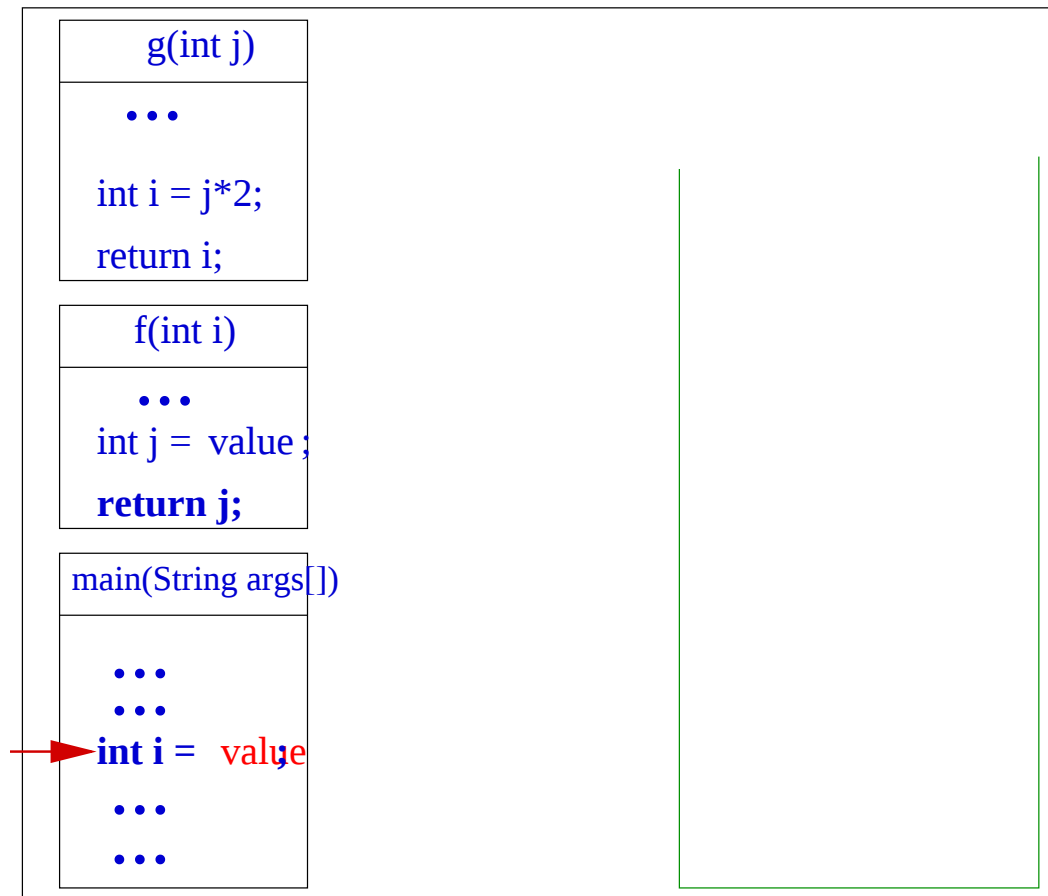
Memory

We explore the top entry of the stack



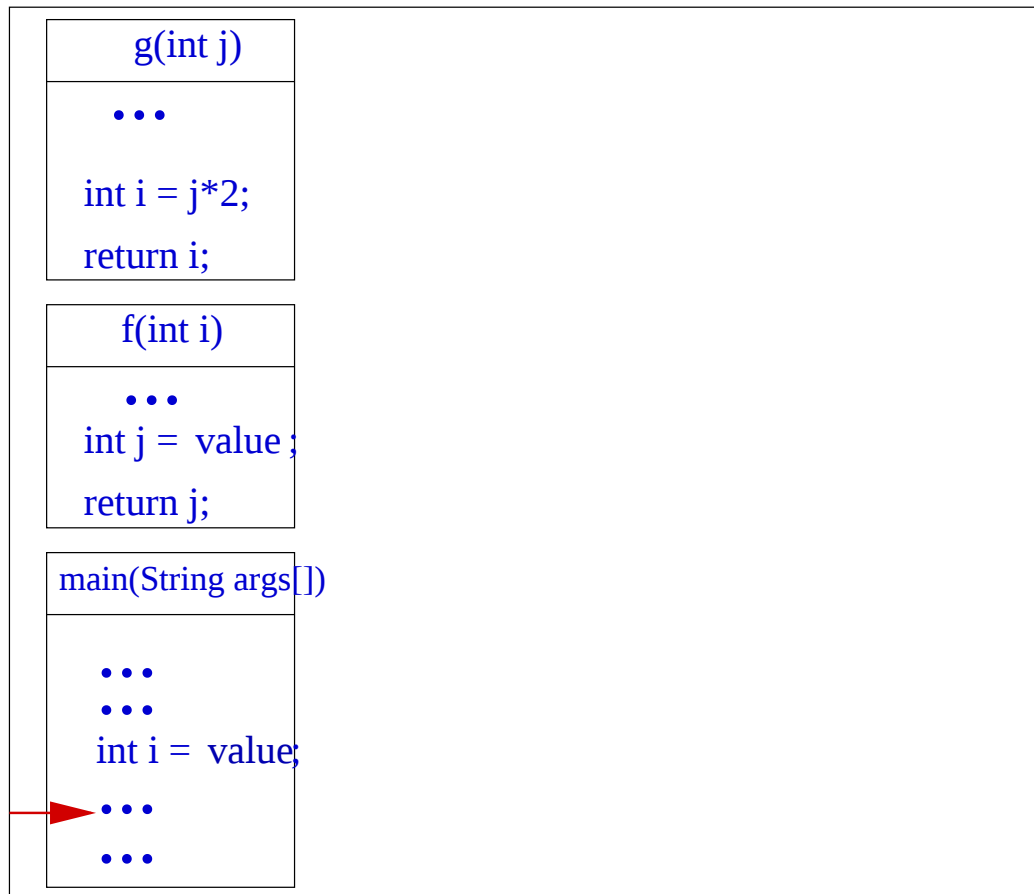
Memory

We copy the value returned by **f ()** to **i**



Memory

The execution proceeds as usual



Memory

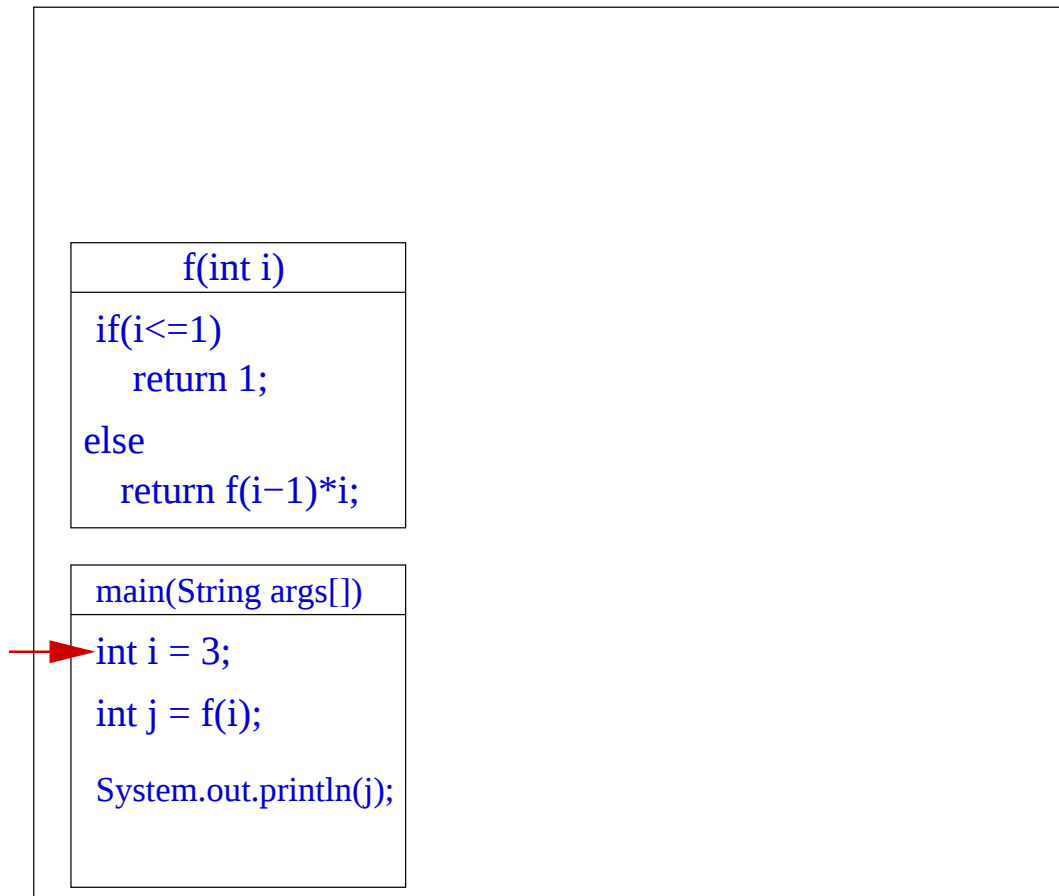
Having got the complete picture of method invocation ...

Let us consider our second program.

Everything is quite systematic now

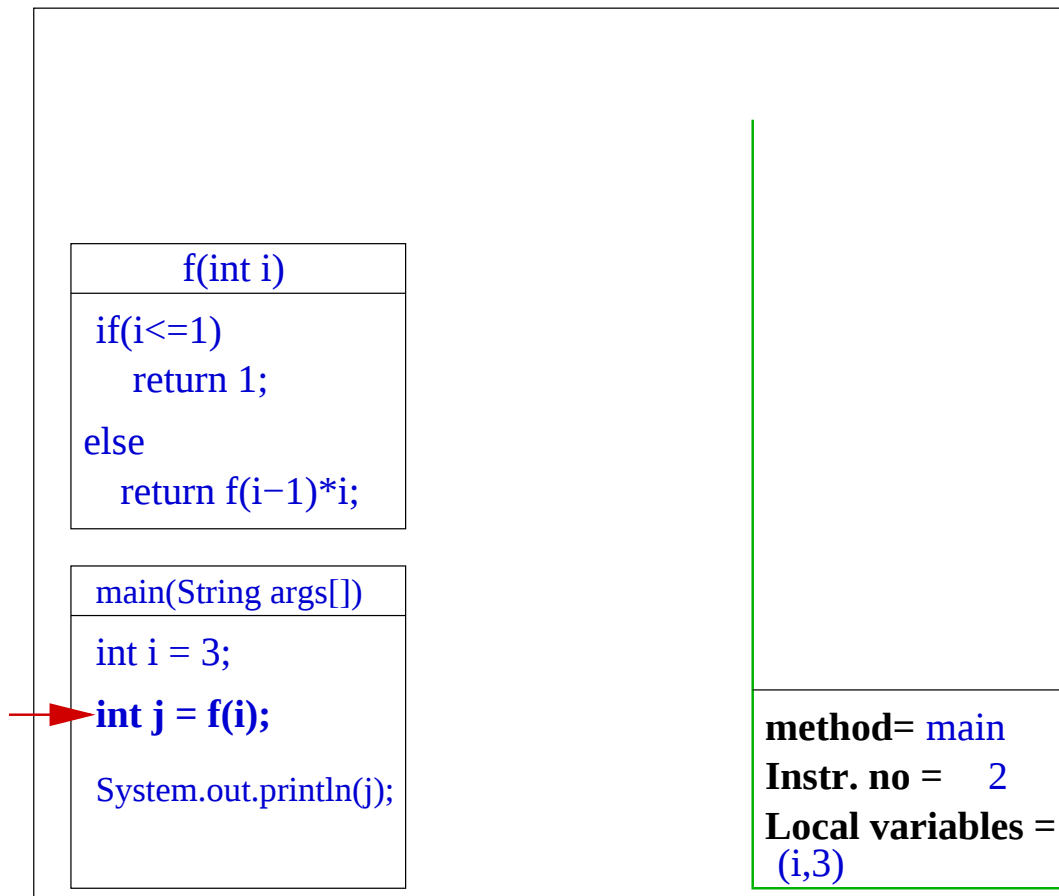
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public static int f(int i)
{ if(i<=1)  return 1;
  else
    return f(i-1)*i;
}
public static void main(String args[])
{  int i = 3;
   int j = f(i);
   System.out.println(j);
}
```

Execution starts from 1st instruction of **main()**



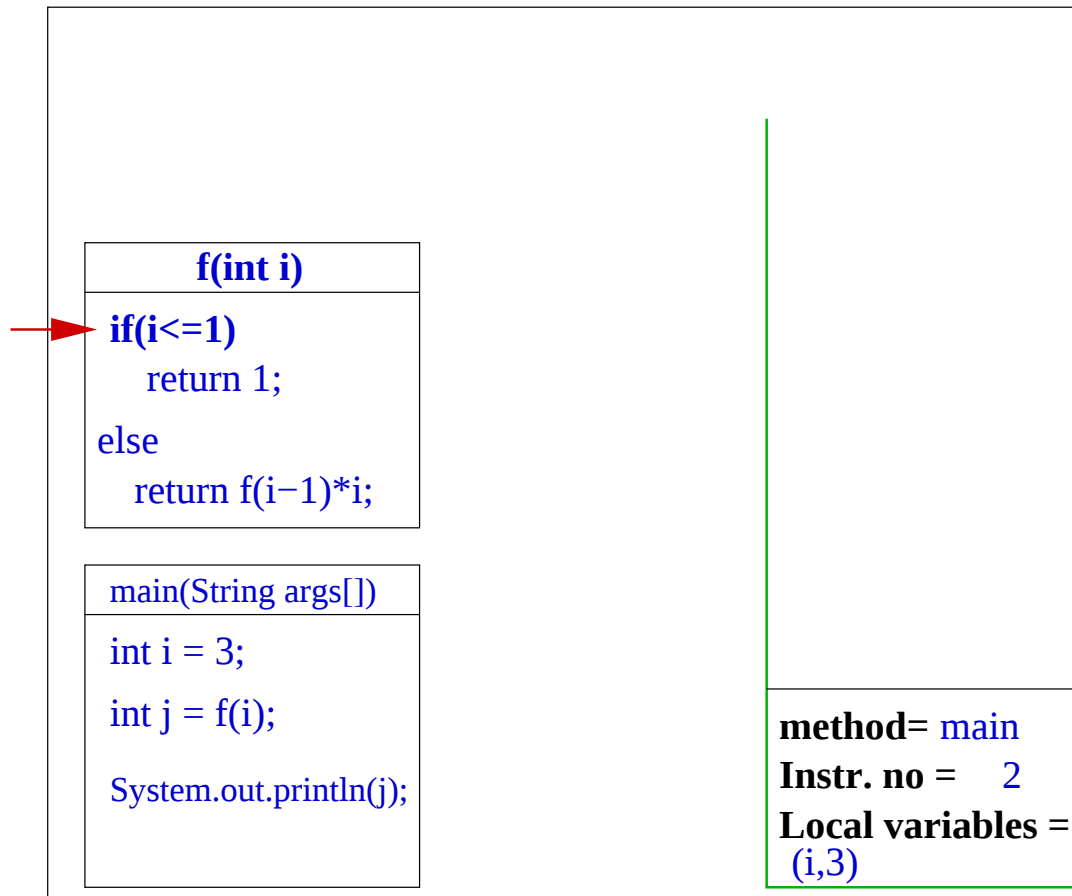
Memory

When we reach instruction `int j = f(i)`



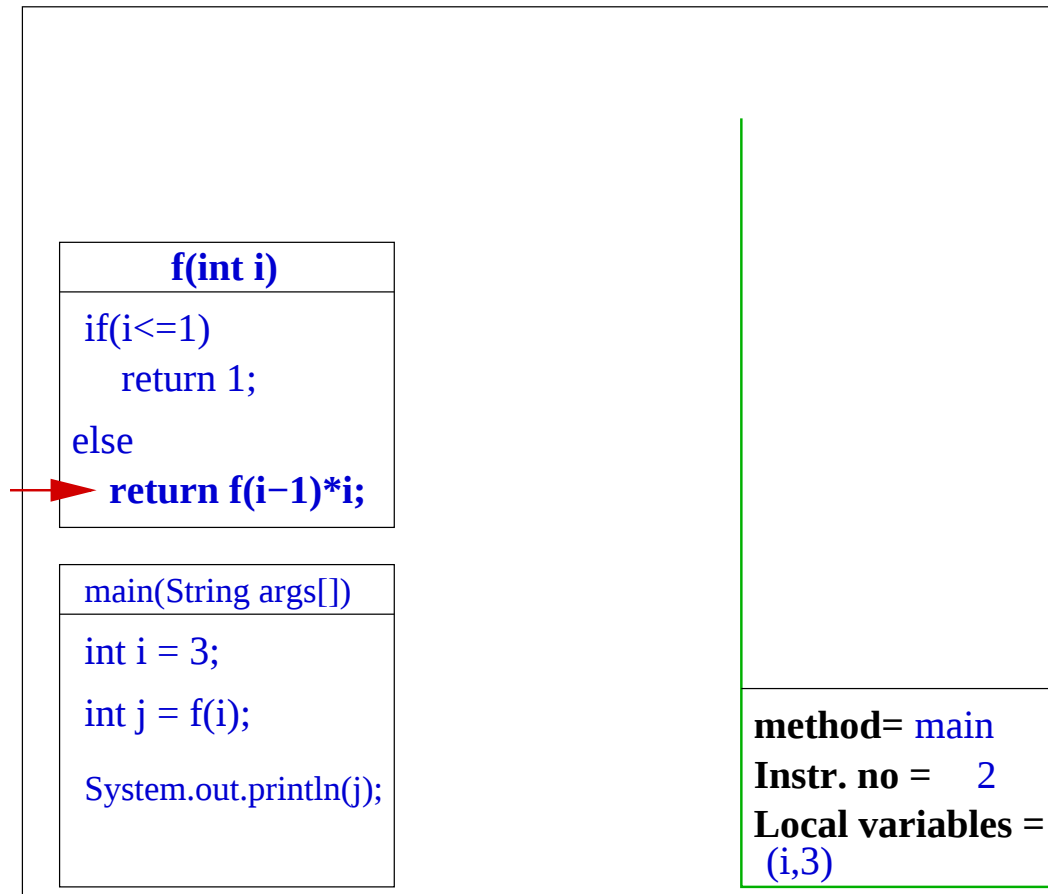
Memory

We keep record of **main** and invoke **f(3)**



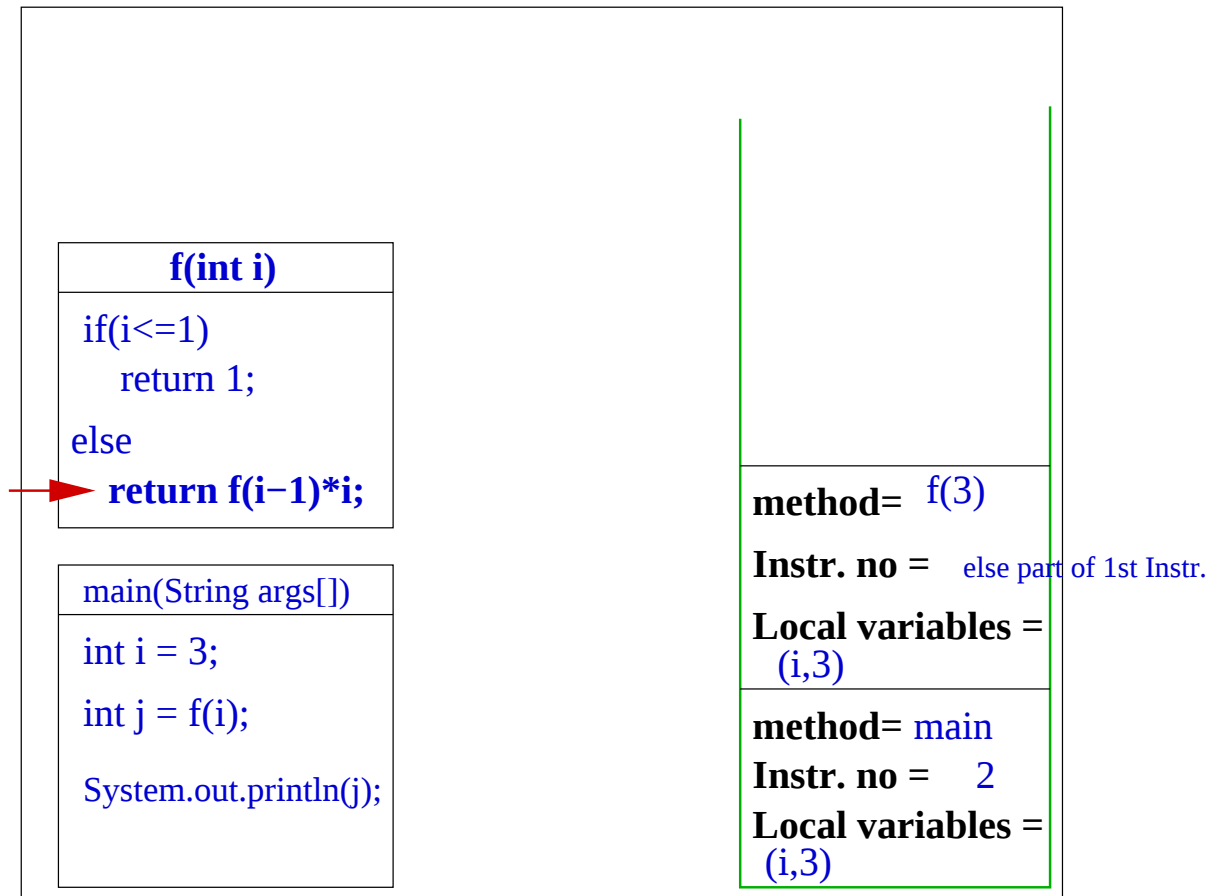
Memory

Since $i > 1$, we reach else part of 1st instr.



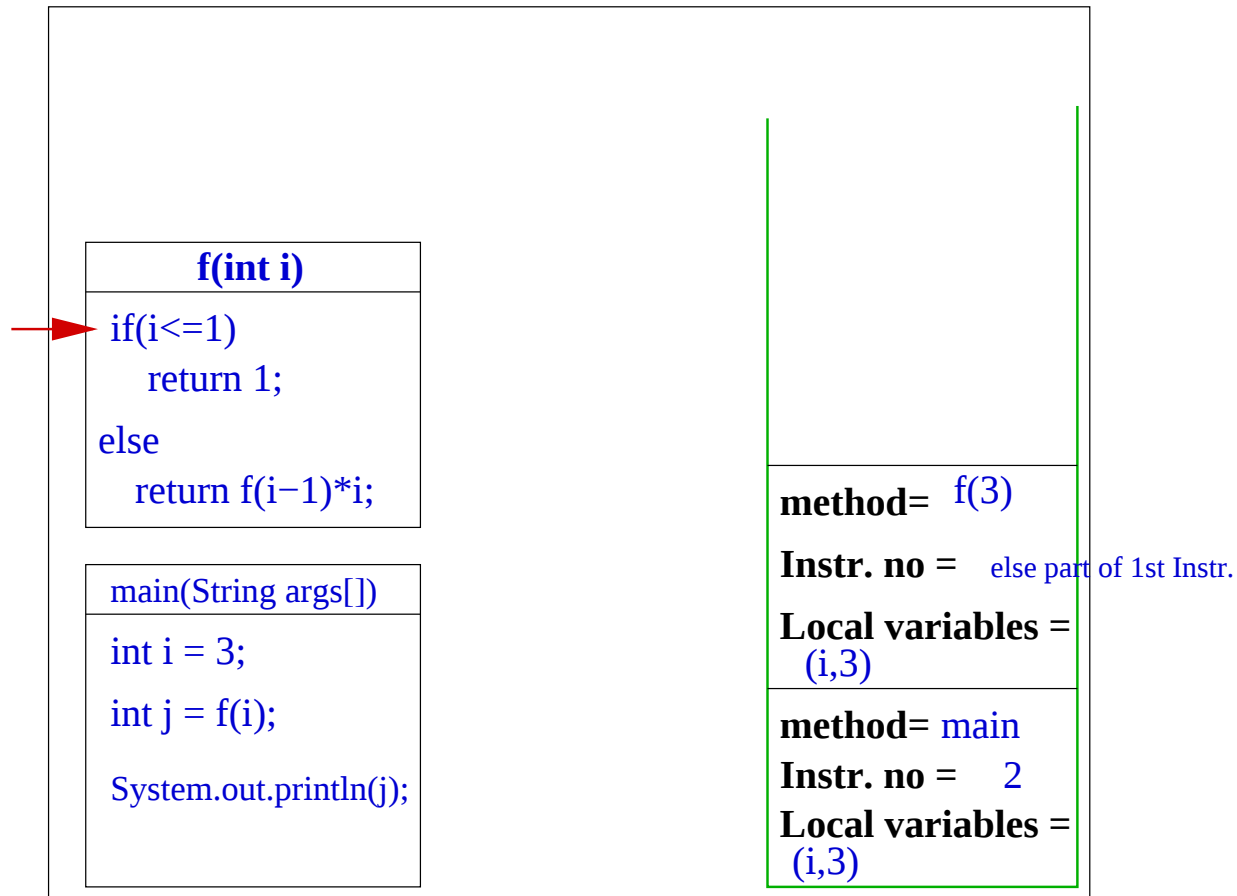
Memory

We keep record of current **f(3)** and invoke **f(2)**



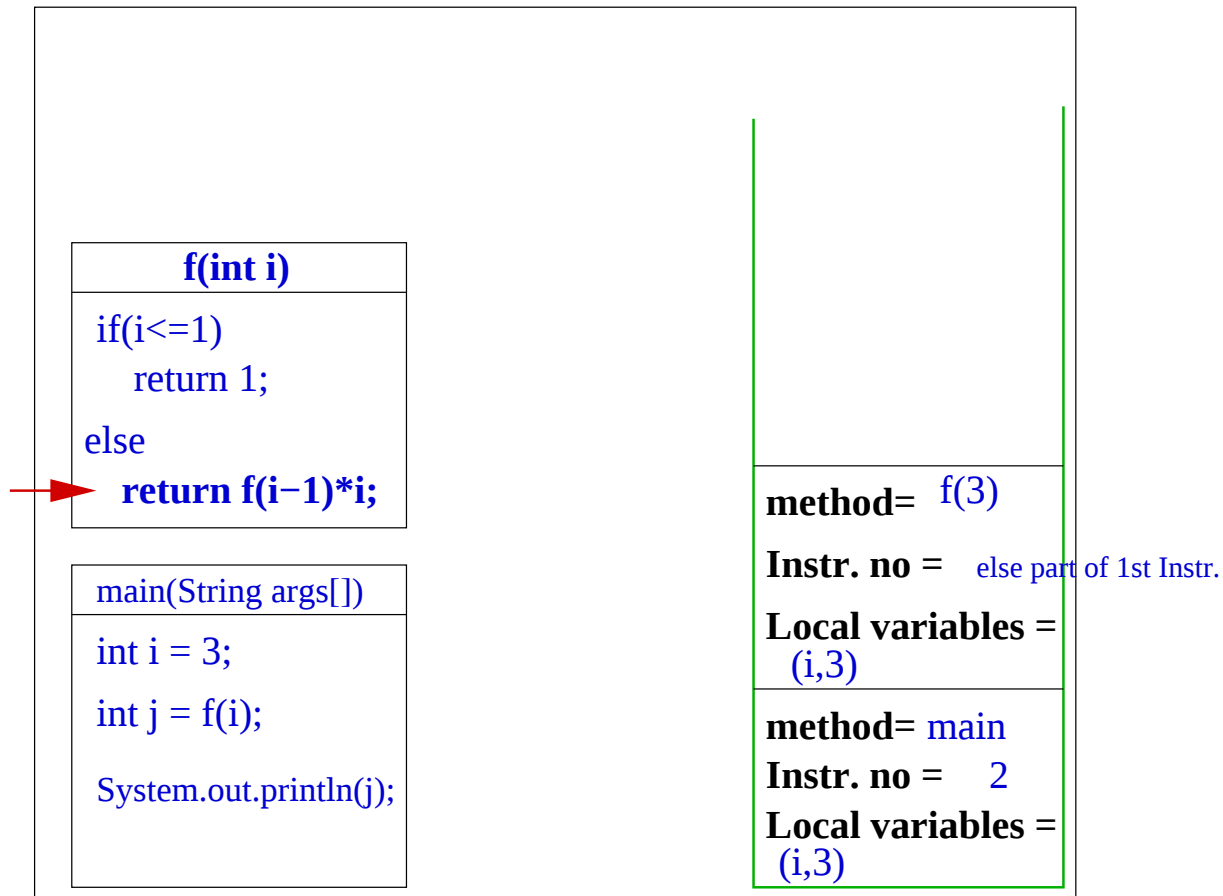
Memory

Since $i > 1$ still, we reach else part of 1st instr.



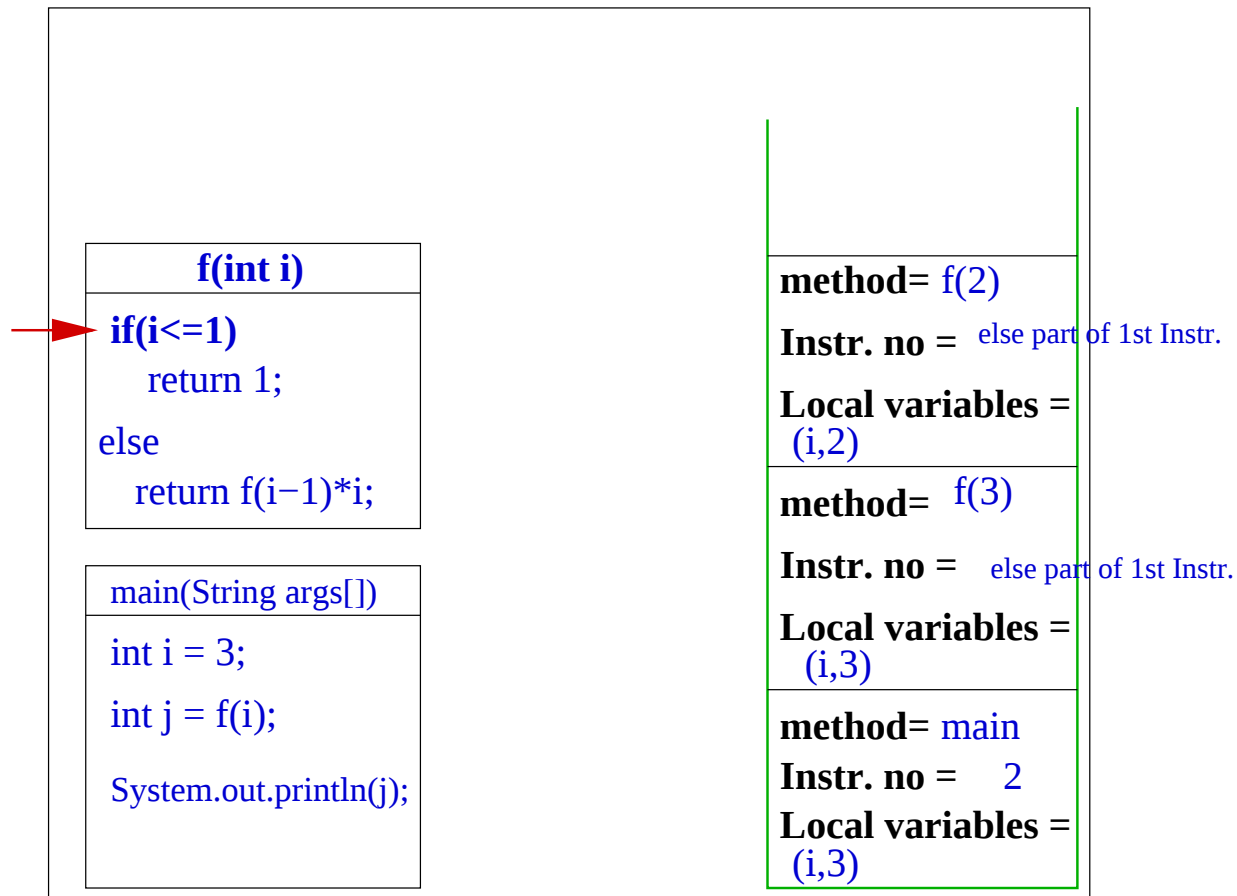
Memory

We keep record of current **f(2)** and invoke **f(1)**



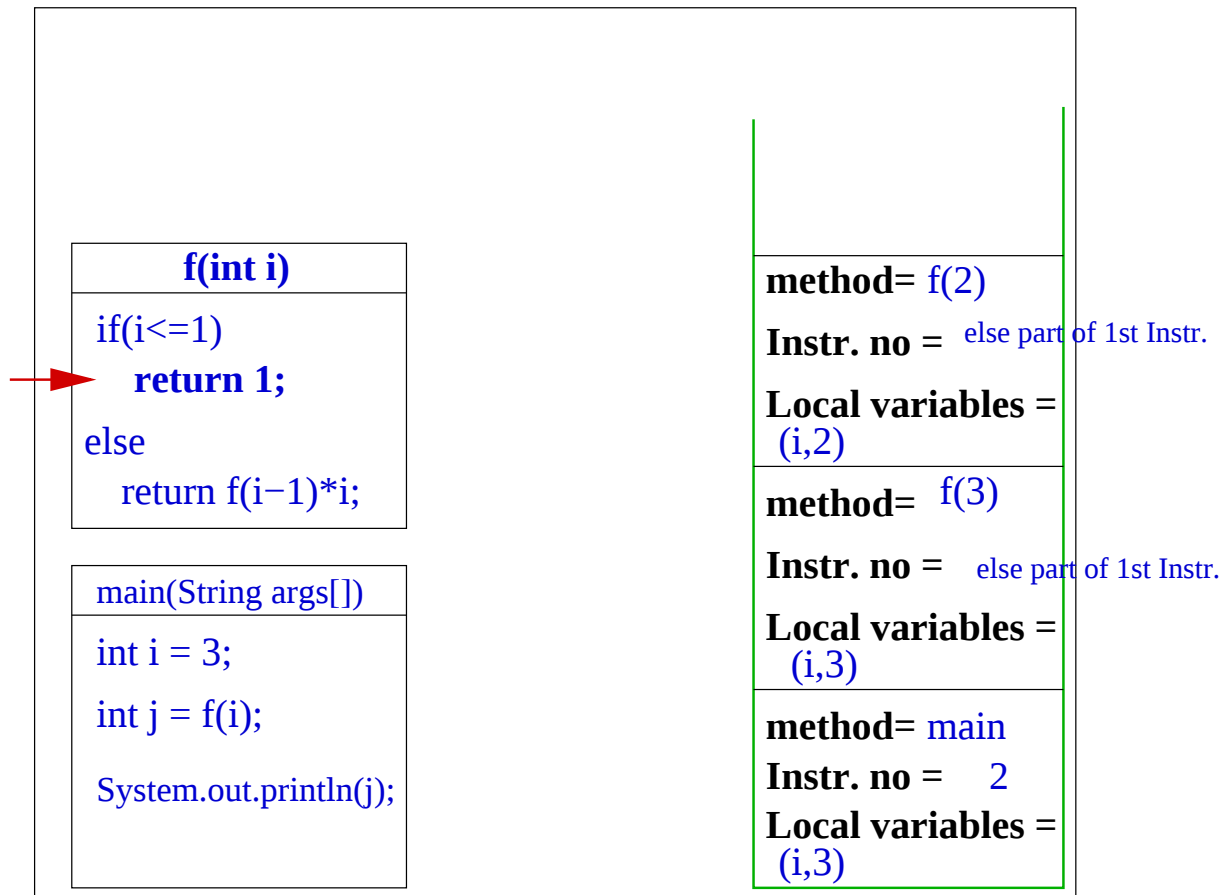
Memory

since $i=1$, we execute the instruction **return 1**



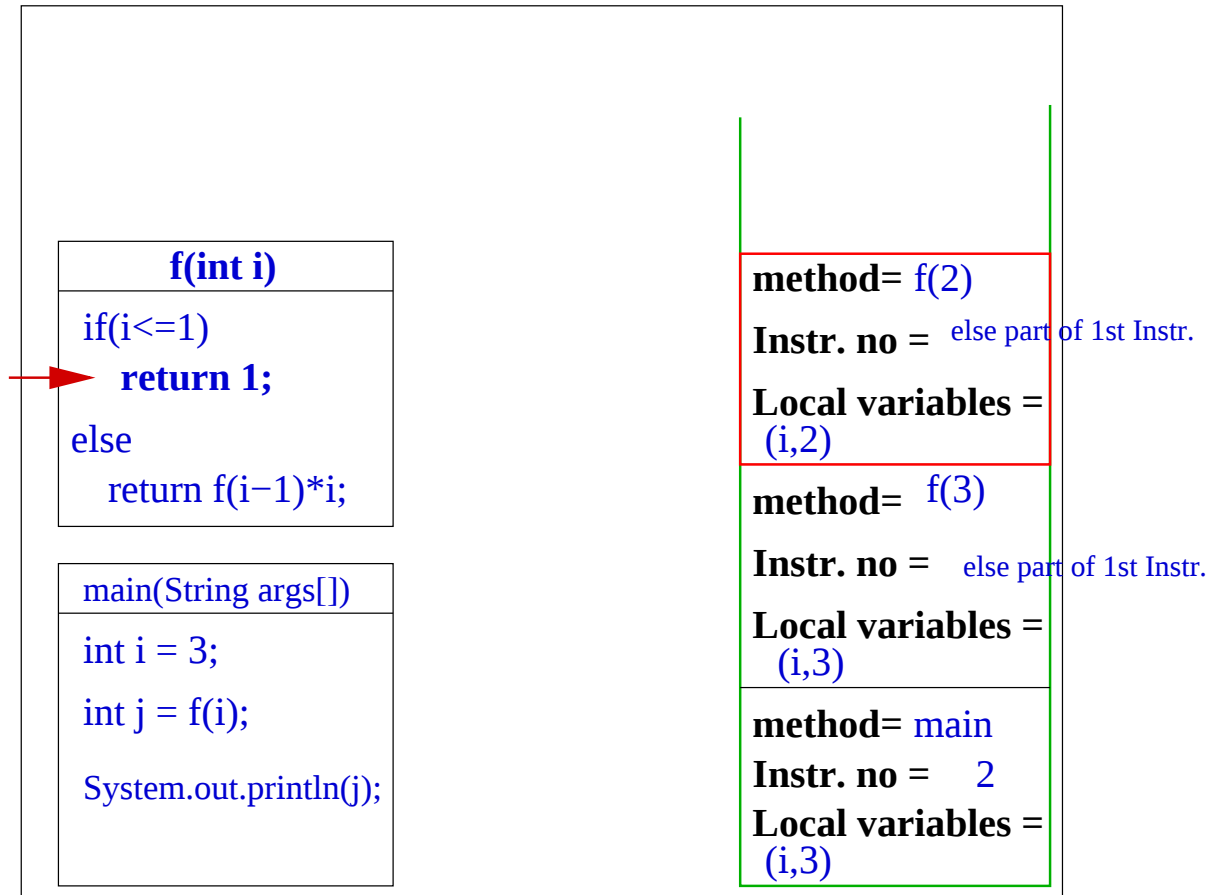
Memory

Where do we return now ?



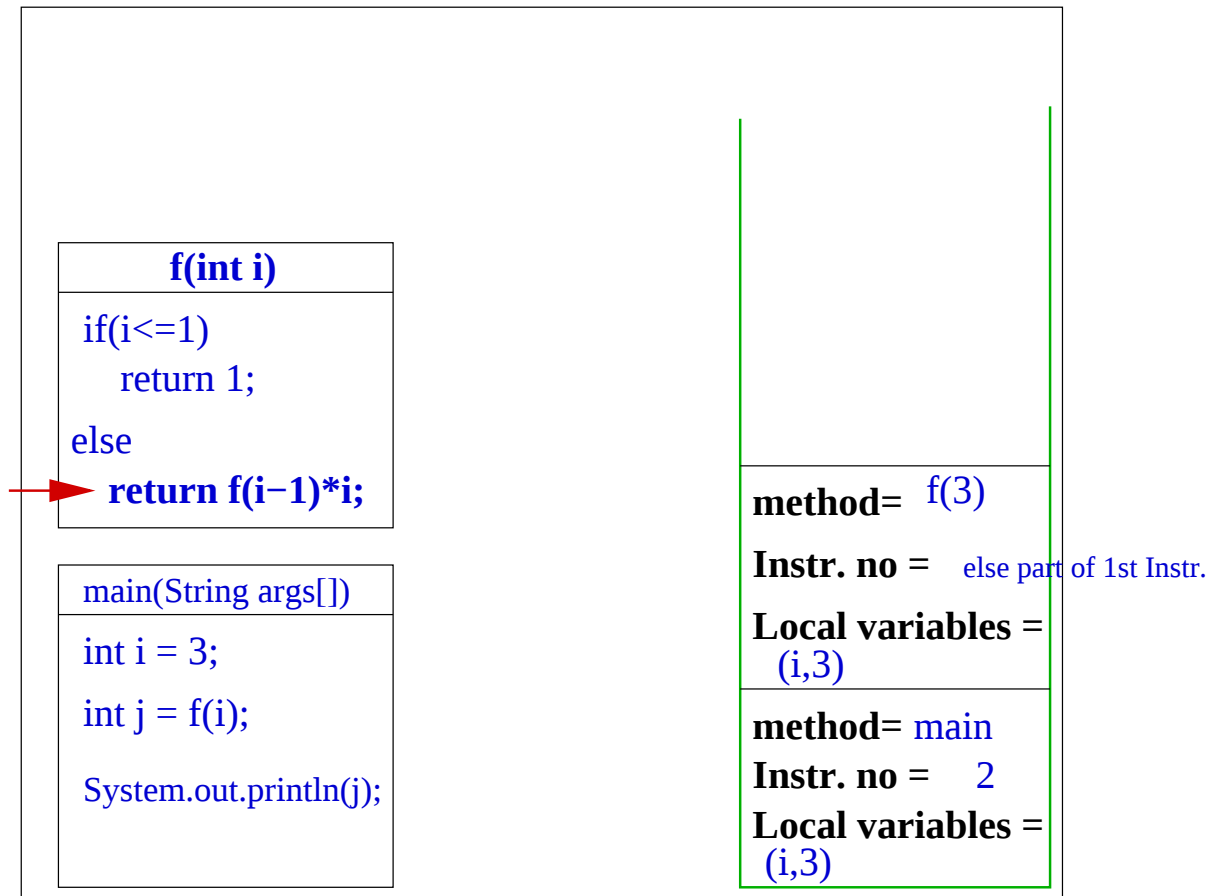
Memory

We return to the method on top of stack : **f(2)** but which instr. ?



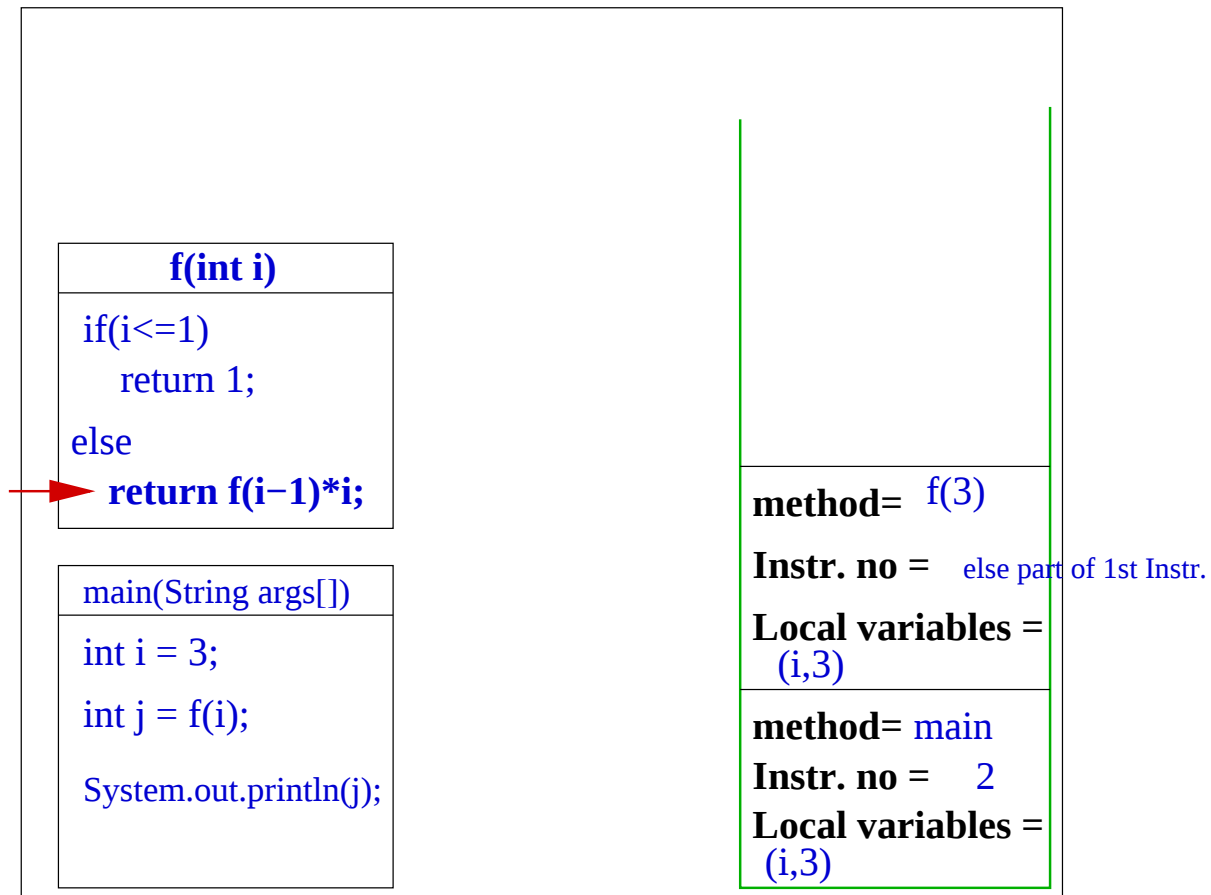
Memory

We return to the else part of the 1st instruction of **f(2)**



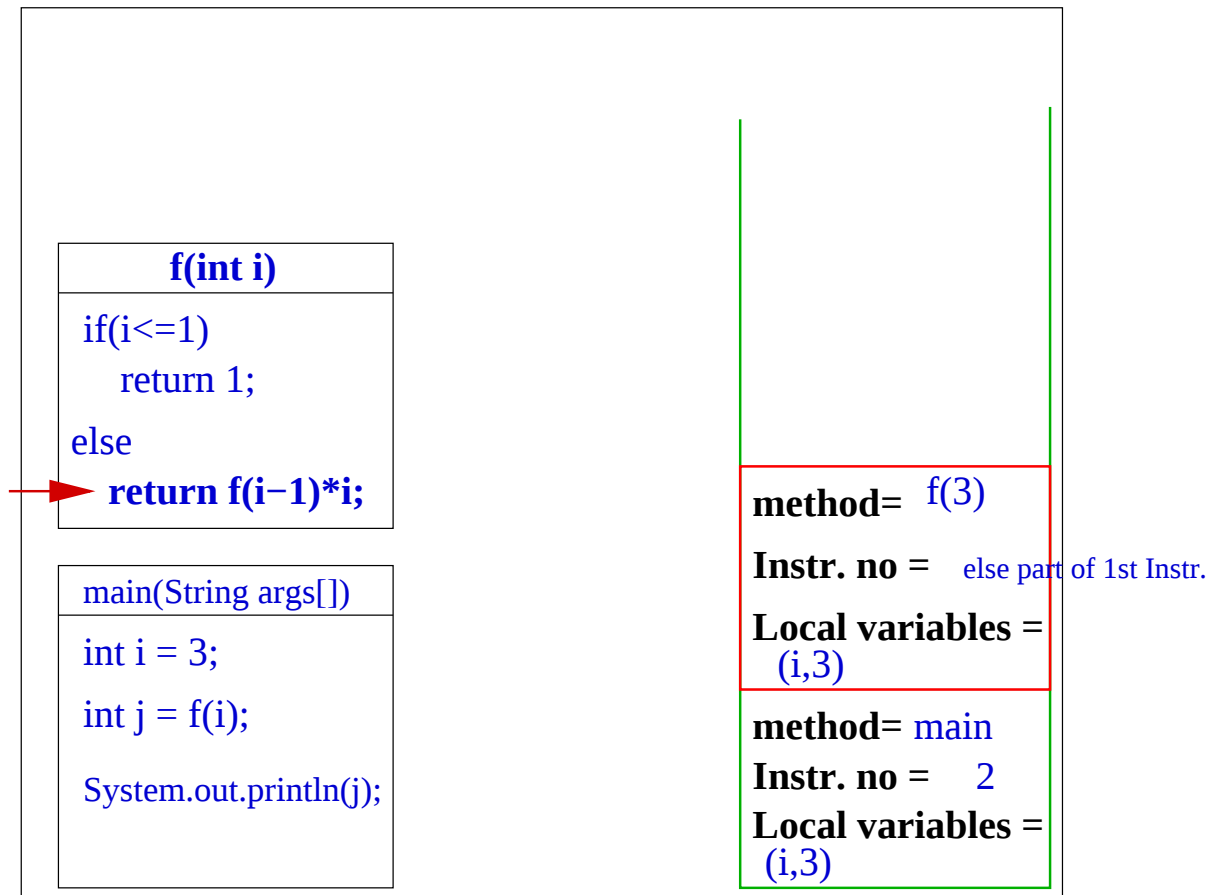
Memory

Where do we return from here ?



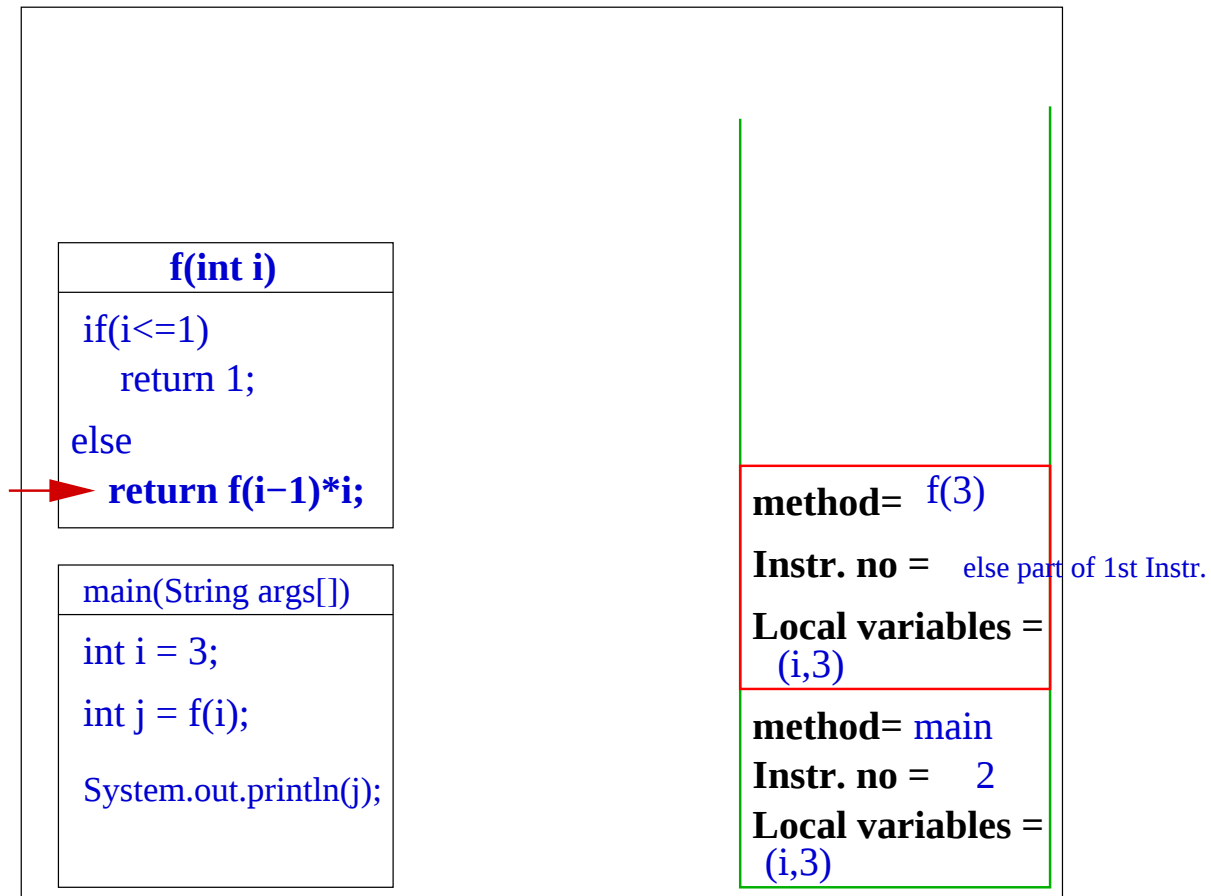
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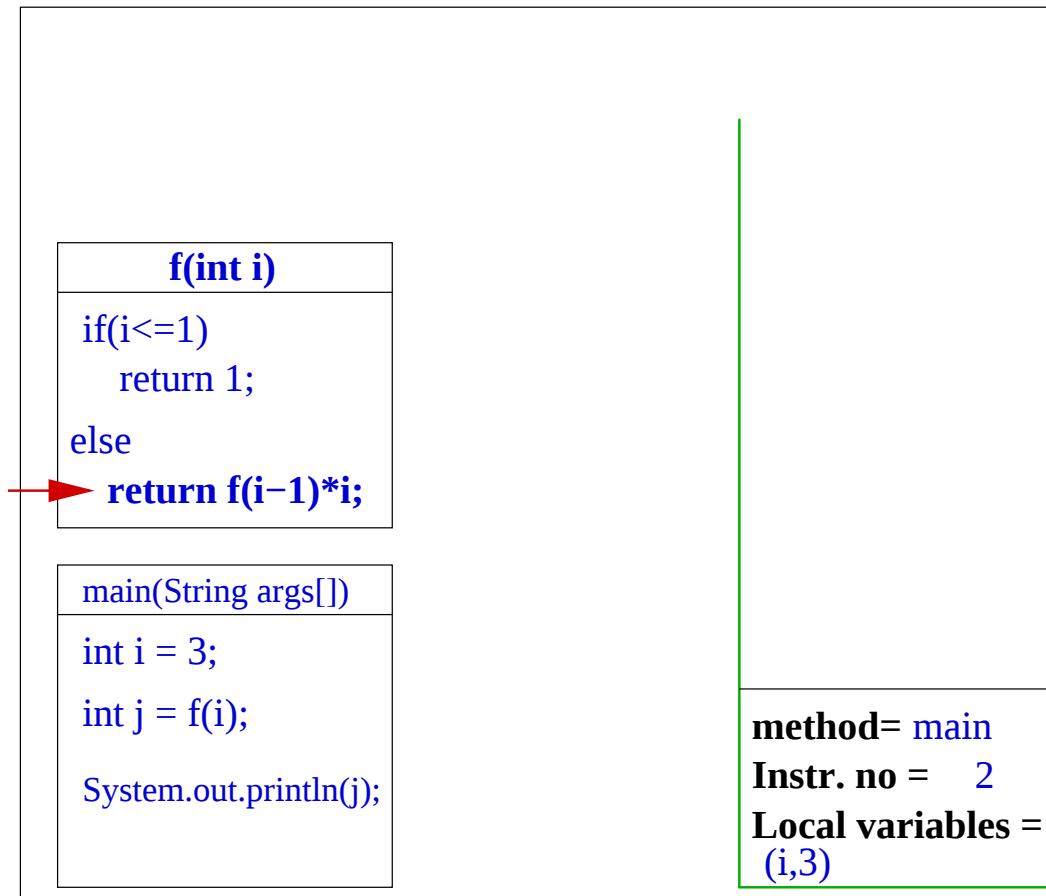
Memory

We return to the method on top of stack : f(3)



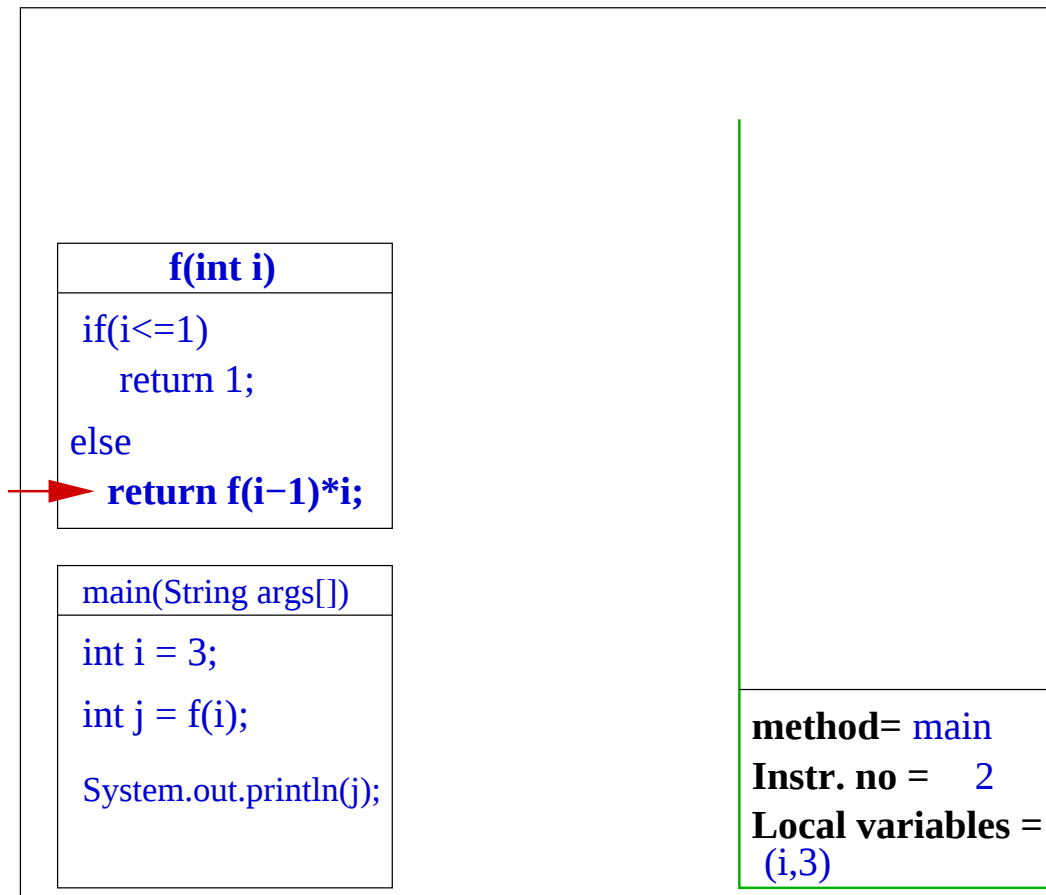
Memory

We are at else part of instruction of method :f(3)



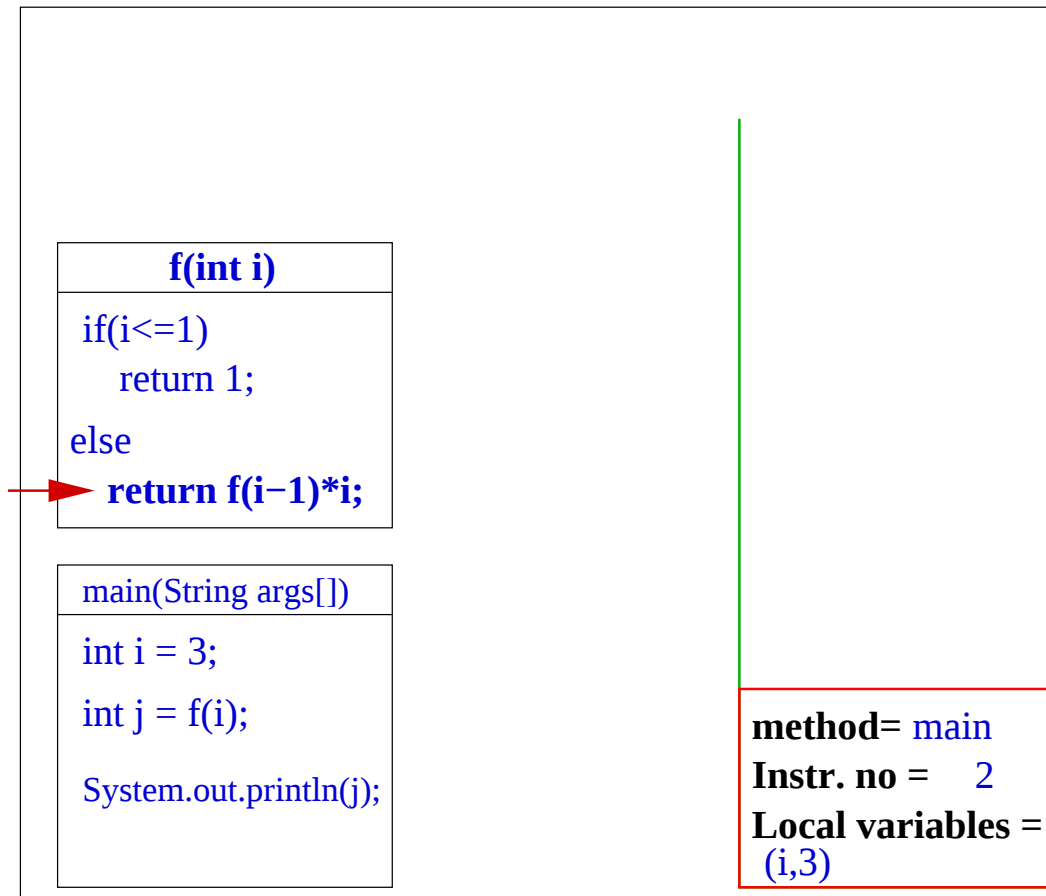
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Where do we return from here ?



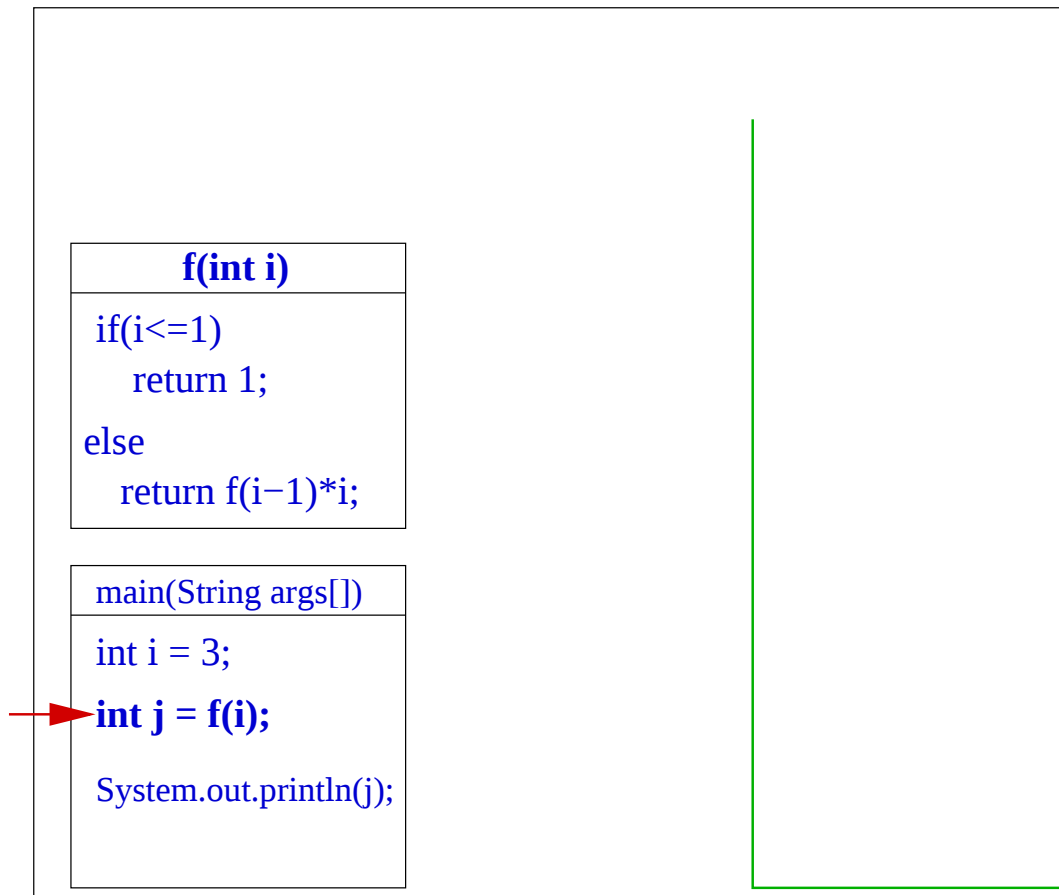
Memory

To the method on top of the stack : **main()**



Memory

and now the usual execution in the main



Memory

Recursion

There are many functions which can be defined inductively :

- fibonacci numbers :

$$f(0) = 0, f(1) = 1;$$

$$\text{for } i > 1 \quad f(i) = f(i - 1) + f(i - 2);$$

- factorial(n) :

$$factorial(0) = 1;$$

$$\text{for } i > 0, \quad factorial(i) = i * factorial(i - 1)$$

Recursion

There are some functions whose definition itself is inductive (recursive) in nature.

- Permutation of string of n characters
- Combinations ...
- Powerset of a set S : the set of all subsets of a the set S

Recursion

There are some functions whose definition itself is inductive (recursive) in nature.

- Permutation of string of n characters
- Combinations ...
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They have a very simple recursive solution