

I Semester 2008-09

Lecture 36

- Announcement : Extra sessions for lab test
- Sorting algorithms based on recursion
 - Quick Sort (did in lst class)
 - Merge Sort
- Introduction to Time Complexity

Extra session for lab

On Sunday- 9th November

B1 and B2	: Nikhi Jain	CS101 2-3:30 PM
B3 and B4	: Kshitiz Garg	CS102 2-3:30 PM.
B5 and B6	: Abhinav/Nitin	CS101 3:45-5:15PM
B7 and B8	: Vikas Mards	CS102 3:45-5:15PM
B9 and B1(: Paras Tikamani	CS101 5:15-6:45PM

Optional extra class

Saturday- 8th November at 5:15 PM in CS101

"Tower of Hanoi" problem

Hints: You can see that for 1 or 2 discs, there is a straight forward solution.

Now for the problem with n + 1 discs, assume you have a black box method which can transfer n discs from one tower to another. How can you solve the given problem (that is, n + 1 discs) using that method.

Now try to fill in the details of the recursive method. The final solution (recursive method) is very small....

Quick Sort

$\label{eq:QuickSort:sortingaset} Soft = Soft a set So$

- Select an element x.
- Partition the set S of numbers into two parts:
 - $S_{<x}$: the subset consisting of numbers less than x.
 - $S_{>x}$: the subset consisting of numbers greater than x.
- Recursively sort the two sets separately $S_{< x}$ and $S_{> x}$, and concatenate them.

Quick Sort : when set is represented as array

```
public static void Qsort(int[] A, int left, int right)
{
    if(left<right)
    {
        int mid = partition(A, left, right);
        Qsort(A,left,mid-1);
        Qsort(A,mid+1,right);
    }
}</pre>
```

You can observe that the size of problem corresponding to recursive calls decreases always. Hence the program will eventually terminate.

Merge Sort

Merging two sorted arrays

Problem : Given two sorted arrays A and B, produce another sorted array $C \bigcup A \cup B$.

Merging two sorted arrays

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Trivial Solution :

Copy the elements of ${\cal A}$ and ${\cal B}$ into ${\cal C},$ then sort ${\cal C}$

Merging two sorted arrays

Problem : Given two sorted arrays A and B, produce another sorted array $C \bigcup A \cup B$.

Trivial Solution :

Copy the elements of A and B into C, then sort CMissing Point : A and B are already sorted

Merging two sorted arrays : a better solution

start scanning A and B from left, compare two elements of A and B, copy the smaller one into C and continue ...

Merging two sorted arrays : a better solution

start scanning A and B from left, compare two elements of A and B, copy the smaller one into C and continue ...

Let us consider an example ...



























Merge Sort

Key Idea : Merging two sorted arrays is easier than sorting their union.

- Sort the first half of array A recursively
- Sort the second half of the array recursively
- merge the two halves.





```
public static void mergesort(int[] A, int left, int right)
        if(left!=right)
            int mid = (left+right)/2;
                       ??
                                      ;
                       ??
                       ??
                                      ;
    }
```

```
public static void mergesort(int[] A, int left, int right)
        if(left!=right)
            int mid = (left+right)/2;
            mergesort(A, left, mid);
                       ??
                                    ;
                       ??
    }
```

```
public static void mergesort(int[] A, int left, int right)
        if(left!=right)
            int mid = (left+right)/2;
            mergesort(A, left, mid);
            mergesort(A, mid+1, right);
                      ??
                                       ;
    }
```

```
public static void mergesort(int[] A, int left, int right)
        if(left!=right)
            int mid = (left+right)/2;
            mergesort(A, left, mid);
            mergesort(A, mid+1, right);
            merge(A,left,mid,right);
```

```
public static void mergesort(int[] A, int left, int right)
{
        if(left!=right)
        {
            int mid = (left+right)/2;
            mergesort(A, left, mid);
            mergesort(A, mid+1, right);
            merge(A,left,mid,right);
        }
}
```

Convince yourself that each recursive call makes progress, that is, it approaches the base case. This is very important rule that each recursive method must obey.

Recursion Tree

When a method makes two or more recursive calls to itself, it is better to view the execution as a tree.

Recursion Tree for Fibonacci number

```
public static int fib(int n)
{
    if(n==0) return 0;
    else
    {       if(n==1) return 1;
            else return fib(n-1)+fib(n-2);
    }
}
Example:n=4
```

Recursion Tree for Fibonacci number for n = 4



Recursion Tree for Merge sort

```
public static void mergesort(int[] A, int left, int right)
         if(left!=right)
             int mid = (left+right)/2;
             mergesort(A, left, mid);
             mergesort(A, mid+1, right);
             merge(A,left,mid,right);
Example : A = \{99, 7, 5, 1, 67, 11, 4, 2\}
```

Recursion Tree for Merge sort

for A = $\{99, 7, 5, 1, 67, 11, 4, 2\}$

Note : In the next few slides, for sake of compactness we shall use MSort() to denote mergesort().

We show the status of A as we move up the recursion tree level by level.















Recursion Tree for Quick sort

Do it as homework



```
Measuring time taken a method \ensuremath{\mathbb{M}}
long start = System.currentTimeMillis();
M();
long stop = System.currentTimeMillis();
System.out.println(stop-start);
```

Note: System.currentTimeMillis() returns a long which corresponds to current time in milliseconds.

Comparing Three sorting algorithms

Experimental Observations

- Quick sort is more efficient than merge sort
- Merge sort is more efficient than Selection Sort

Please study the program : Three_sorting_algos.java

What is the reason for different running times ?

Given that

- all of them has same input and output
- all of them are executed on the same machine

We need to analyze the number of steps/instruction taken by each sorting algorithm

```
for(int i=1; i<=n; i=i+1)
{</pre>
```

```
sum = sum + i;
```

Steps =

```
for(int i=1; i<=n; i=i+1)
{</pre>
```

```
sum = sum + i;
```

```
Steps = 1 + 3n + 1
```

```
for(int n=1;n<=m;n=n+1)
{
    for(int i=1; i<=n; i=i+1)
    {
        sum = sum + i;
    }
}
Steps =</pre>
```

```
for(int n=1;n<=m;n=n+1)
{
    for(int i=1; i<=n; i=i+1)
    {
        sum = sum + i;
    }
}
Steps = 1 + m + \sum_{n=1}^{n=m} (1 + 3n + 1) + m + 1
```